



Arrangements and Combinatorics

Richard P. Stanley

M.I.T.

The main reference

An introduction to hyperplane arrangements,
in *Geometric Combinatorics* (E. Miller, V. Reiner,
and B. Sturmfels, eds.), IAS/Park City
Mathematics Series, vol. 13, American
Mathematical Society, Providence, RI, 2007,
pp. 389–496.

math.mit.edu/~rstan/arrangements/arr.html

Posets

A **poset** (partially ordered set) is a set P and relation \leq satisfying $\forall x, y, z \in P$:

(P1) (**reflexivity**) $x \leq x$

(P2) (**antisymmetry**) If $x \leq y$ and $y \leq x$, then $x = y$.

(P3) (**transitivity**) If $x \leq y$ and $y \leq z$, then $x \leq z$.

Arrangements

K : a field

\mathcal{A} : a (finite) arrangement in $V = K^n$

$\text{rk}(\mathcal{A})$ (**rank** of \mathcal{A}) : dimension of space
spanned by normals to $H \in \mathcal{A}$

Subspaces X, Y, W

Y = any complement to subspace X of K^n
spanned by normals to $H \in \mathcal{A}$

$$W = \{v \in V : v \cdot y = 0 \ \forall y \in Y\}.$$

If $\text{char}(K) = 0$ can take $W = X$.

Essentialization

$$\operatorname{codim}_W(H \cap W) = 1, \quad \forall H \in \mathcal{A}$$

Essentialization of \mathcal{A} :

$$\text{ess}(\mathcal{A}) = \{H \cap W : H \in \mathcal{A}\},$$

an arrangement in W .

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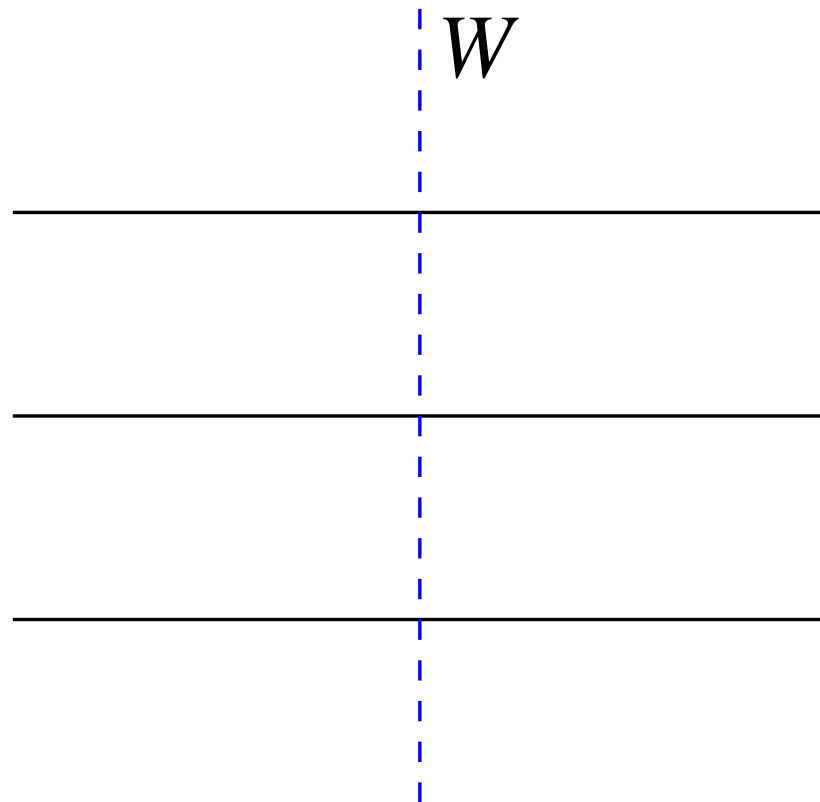
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an arrangement in W .

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\mathcal{A} is **essential** if $\text{ess}(\mathcal{A}) = \mathcal{A}$, i.e.,
 $\operatorname{rk}(\mathcal{A}) = \dim(\mathcal{A})$.

Example of essentialization



A



$\text{ess}(A)$

The intersection poset

$L(\mathcal{A})$: **nonempty** intersections of hyperplanes
in \mathcal{A} , ordered by **reverse** inclusion

Include V as the bottom element of $L(\mathcal{A})$,
denoted $\hat{0}$.

Note. $L(\mathcal{A}) \cong L(\text{ess}(\mathcal{A}))$

The intersection poset

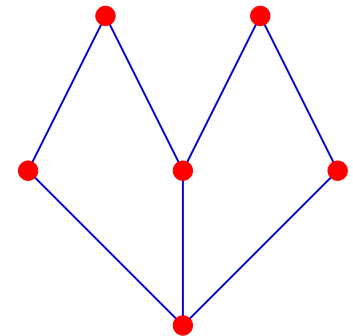
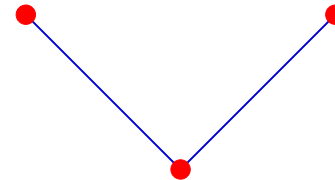
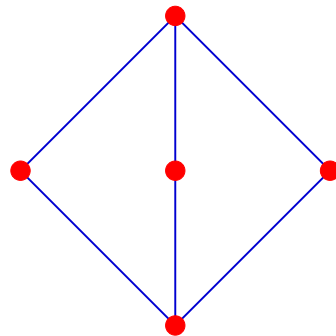
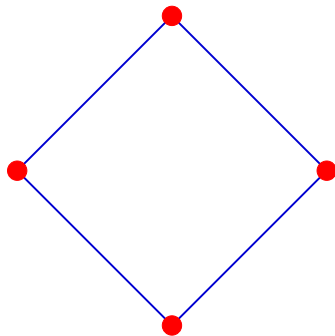
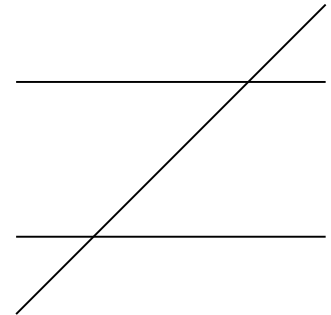
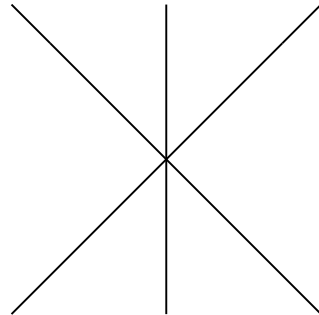
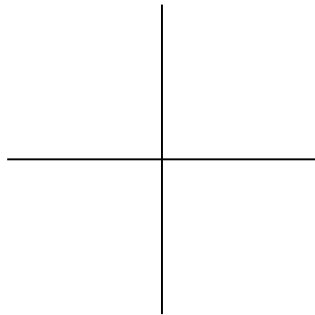
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$L(\mathcal{A})$ is the most important combinatorial object
associated with \mathcal{A} .

Examples of intersection posets



Rank function

Chain of length k : $x_0 < x_1 < \cdots < x_k$

Graded poset of rank n : every maximal chain has length n

Rank function: $\rho(x)$ is the length k of longest chain $x_0 < x_1 < \cdots < x_k = x$.

Rank function on $L(\mathcal{A})$

Proposition. $L(\mathcal{A})$ is graded of rank equal to $\text{rk}(\mathcal{A})$. Rank function:

$$\text{rk}(x) = \text{codim}(x) = n - \dim(x),$$

where $\dim(x)$ is the dimension of x as an affine subspace of V .

Rank function on $L(\mathcal{A})$

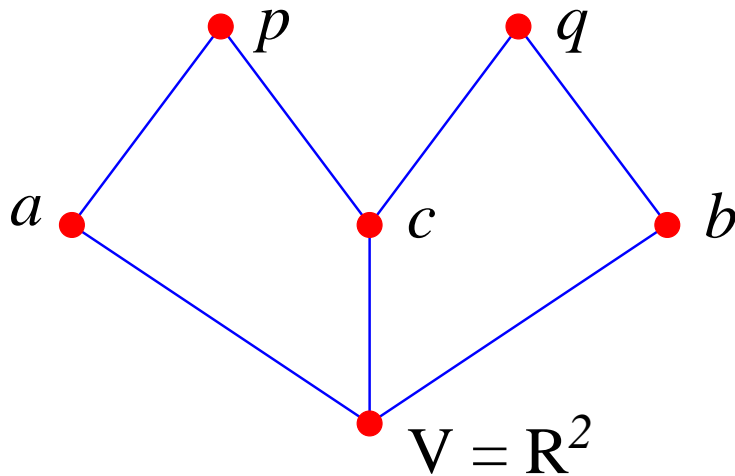
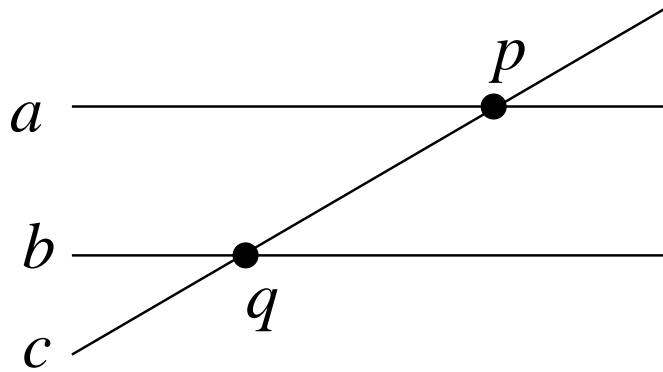
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Proof. Straightforward. \square

Example of $L(\mathcal{A})$



<i>rank</i>	<i>dim</i>
2	0
1	1
0	2

The Möbius function

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Define $\mu = \mu_P : \text{Int}(P) \rightarrow \mathbb{Z}$ (the **Möbius function** of P) by:

$$\mu(x, x) = 1, \text{ for all } x \in P$$

$$\mu(x, y) = - \sum_{x \leq z < y} \mu(x, z), \text{ for all } x < y \text{ in } P.$$

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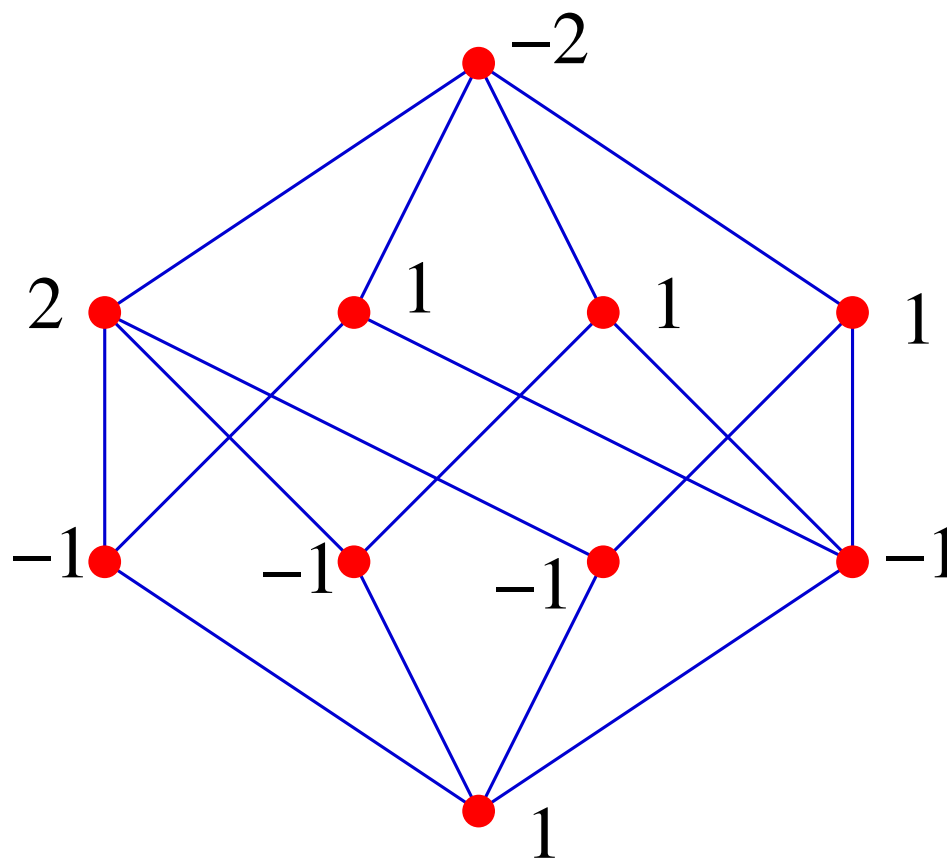
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Write $\mu(x) = \mu(\hat{0}, x)$.

Example of Möbius function



Numbers denote $\mu(x)$.

Möbius inversion formula

P = finite poset

$f, g: P \rightarrow L$ (a field, or even just an abelian group)

Theorem. *Equivalent:*

$$f(x) = \sum_{y \geq x} g(y), \text{ for all } x \in P$$

$$g(x) = \sum_{y \geq x} \mu(x, y) f(y), \text{ for all } x \in P.$$

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The characteristic polynomial

Definition. The *characteristic polynomial* $\chi_{\mathcal{A}}(t)$ of the arrangement \mathcal{A} is defined by

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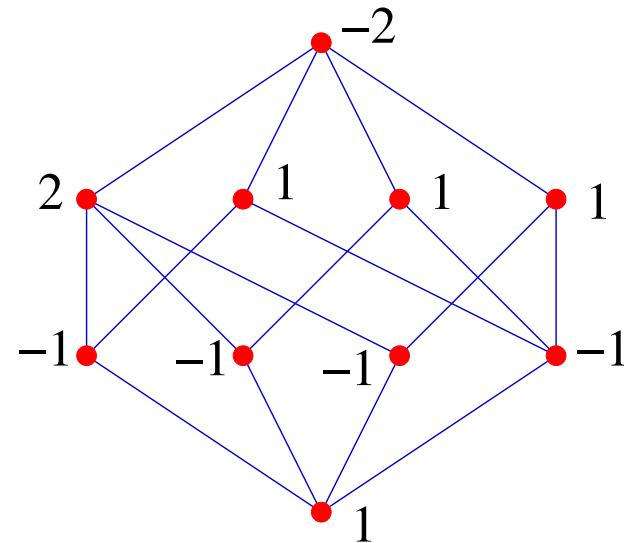
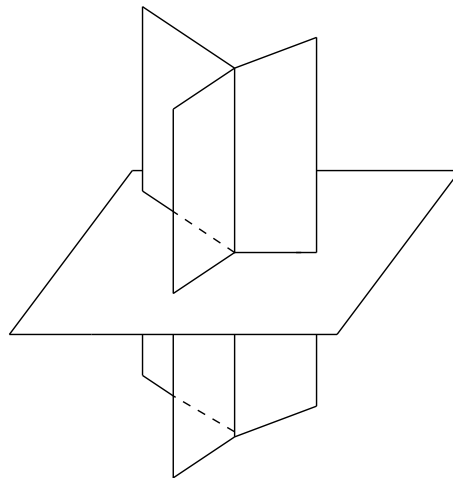
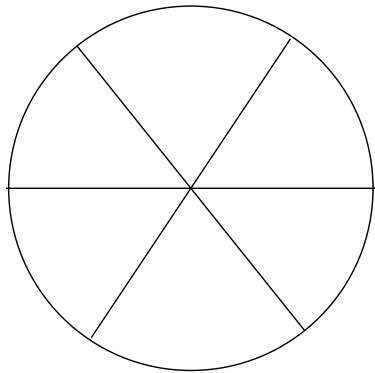
$$\chi_{\mathcal{A}}(t) = \sum_{x \in L(\mathcal{A})} \mu(x) t^{\dim(x)}.$$

Note. $x = V$ contributes t^n , and each $H \in \mathcal{A}$ contributes $-t^{n-1}$. Hence

$$\chi_{\mathcal{A}}(t) = t^n - (\#\mathcal{A})t^{n-1} + \dots.$$

An example

Example.



$$\chi_{\mathcal{A}}(t) = t^3 - 4t^2 + 5t - 2 = (t - 1)^2(t - 2).$$

The boolean algebra

Suppose all hyperplanes in \mathcal{A} are linearly independent, and $\#\mathcal{A} = n$. Then all intersections are nonempty and distinct, so

$$L(\mathcal{A}) \cong B_n,$$

the **boolean algebra** of all subsets of $[n] = \{1, \dots, n\}$, ordered by inclusion.

Characteristic polynomial of B_n

Easy induction argument: $\mu(\hat{0}, x) = (-1)^{n - \dim x}$.
Hence

$$\chi_{\mathcal{A}}(t) = \sum_{i=0}^n \binom{n}{i} (-1)^{n-i} t^i = (t - 1)^n.$$

Regions

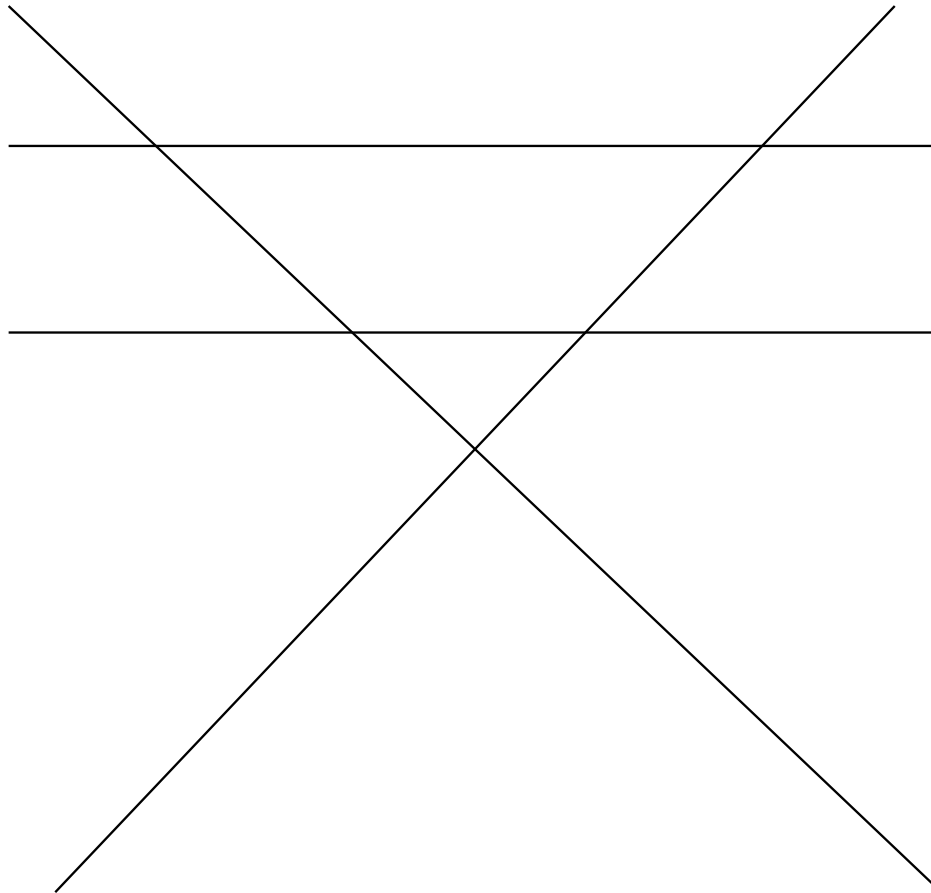
Let $K = \mathbb{R}$. **Region** (or **chamber**) of \mathcal{A} :
connected component of $\mathbb{R}^n - \bigcup_{H \in \mathcal{A}} H$.

$r(\mathcal{A})$ = number of regions of \mathcal{A}

A region R of \mathcal{A} is **relatively bounded** if it becomes bounded in $\text{ess}(\mathcal{A})$.

$b(\mathcal{A})$ = number of relatively bounded regions of \mathcal{A}

Example of $r(\mathcal{A})$ and $b(\mathcal{A})$



$$r(\mathcal{A}) = 10, \quad b(\mathcal{A}) = 2$$

Zaslavsky's theorem (1975)

Current goal:

Theorem. *Let \mathcal{A} be an arrangement of rank r in \mathbb{R}^n . Then*

$$r(\mathcal{A}) = (-1)^n \chi_{\mathcal{A}}(-1)$$

$$b(\mathcal{A}) = (-1)^r \chi_{\mathcal{A}}(1).$$

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Proof will be by induction on $\#\mathcal{A}$ (the number of regions).

Subarrangements and restrictions

subarrangement of \mathcal{A} : a subset $\mathcal{B} \subseteq \mathcal{A}$

For $x \in L(\mathcal{A})$ define

$$\mathcal{A}_x = \{H \in \mathcal{A} : x \subseteq H\} \subseteq \mathcal{A}$$

Subarrangements and restrictions

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For $x \in L(\mathcal{A})$ define

$$\mathcal{A}_x = \{H \in \mathcal{A} : x \subseteq H\} \subseteq \mathcal{A}$$

Also define the **restriction** of \mathcal{A} to x to be the arrangement in the affine space \mathcal{A} :

$$\mathcal{A}^x = \{x \cap H \neq \emptyset : H \in \mathcal{A} - \mathcal{A}_x\}.$$

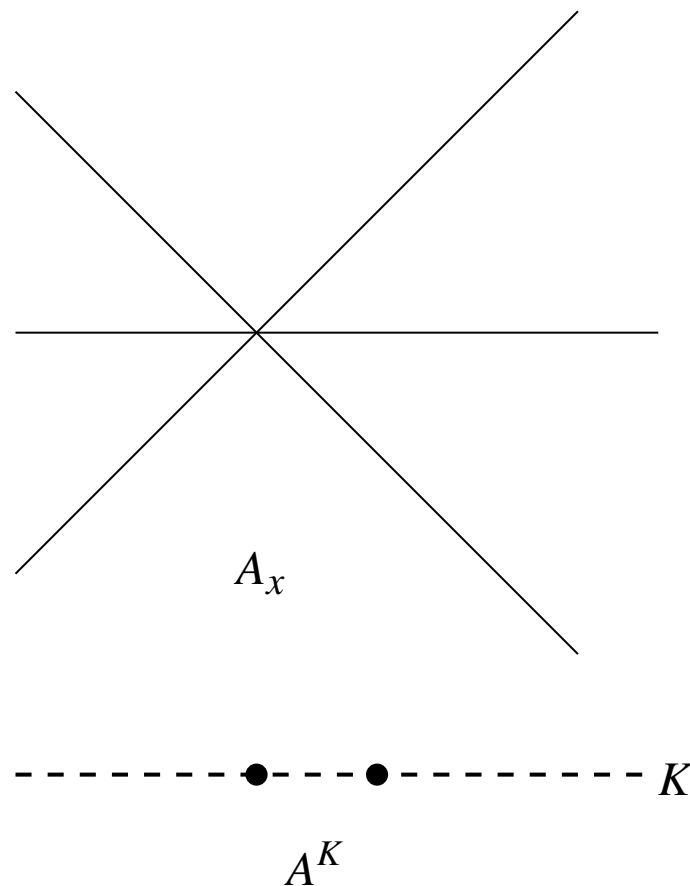
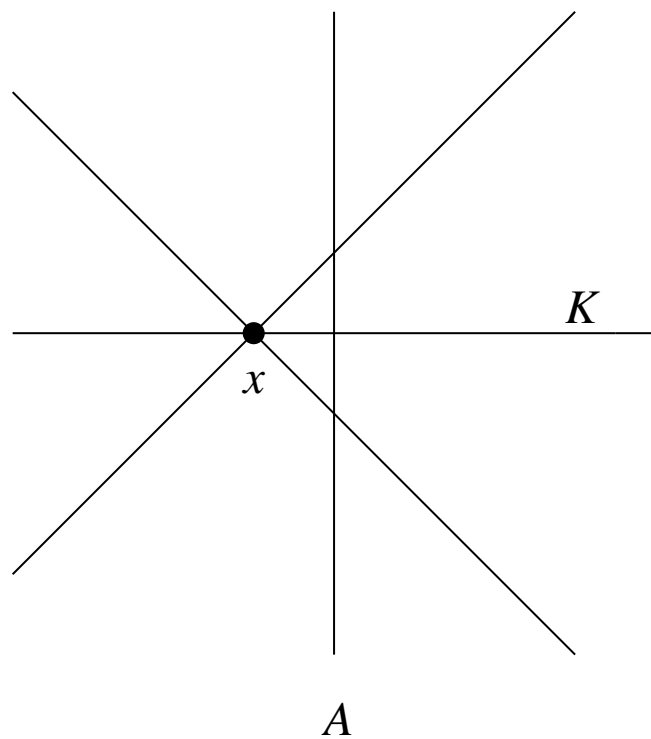
$L(\mathcal{A}_x)$ and $L(\mathcal{A}^x)$

Note that if $x \in L(\mathcal{A})$, then

$$L(\mathcal{A}_x) \cong \Lambda_x := \{y \in L(\mathcal{A}) : y \leq x\}$$

$$L(\mathcal{A}^x) \cong V_x := \{y \in L(\mathcal{A}) : y \geq x\}.$$

Example of \mathcal{A}_x and \mathcal{A}^x



Triple of arrangements

Choose $H_0 \in \mathcal{A}$. Define

$$\mathcal{A}' = \mathcal{A} - \{H_0\}$$

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Call $(\mathcal{A}, \mathcal{A}', \mathcal{A}'')$ a **triple** of arrangements with **distinguished hyperplane** H_0 .

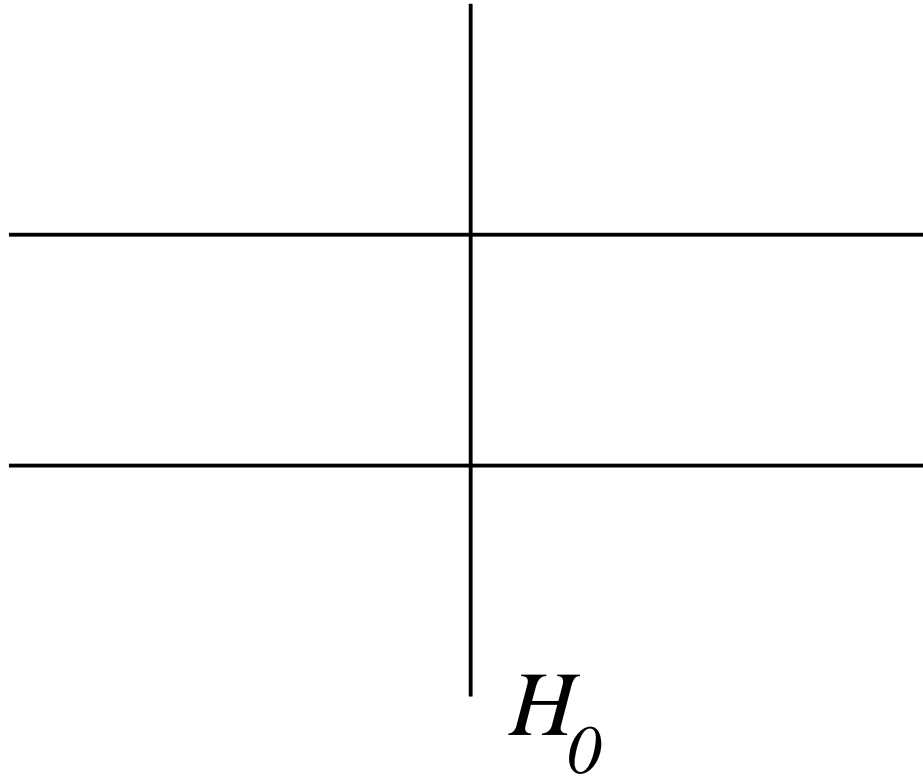
Recurrence for $r(\mathcal{A})$ and $b(\mathcal{A})$

Lemma. *Let $(\mathcal{A}, \mathcal{A}', \mathcal{A}'')$ be a triple of real arrangements with distinguished hyperplane H_0 . Then*

$$r(\mathcal{A}) = r(\mathcal{A}') + r(\mathcal{A}'')$$

$$b(\mathcal{A}) = \begin{cases} b(\mathcal{A}') + b(\mathcal{A}''), & \text{if } \text{rk}(\mathcal{A}) = \text{rk}(\mathcal{A}') \\ 0, & \text{if } \text{rk}(\mathcal{A}) = \text{rk}(\mathcal{A}') + 1. \end{cases}$$

The case $\text{rk}(\mathcal{A}) = \text{rk}(\mathcal{A}') + 1$



Proof of lemma (sketch)

Note that $r(\mathcal{A})$ equals $r(\mathcal{A}')$ plus the number of regions of \mathcal{A}' cut into two regions by H_0 . Easy to give a bijection between regions of \mathcal{A}' cut in two by H_0 and regions of \mathcal{A}'' , proving

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Proof of recurrence for $b(\mathcal{A})$ analogous. \square

The deletion-restriction recurrence

Lemma. Let $(\mathcal{A}, \mathcal{A}', \mathcal{A}'')$ be a triple of real arrangements. Then

$$\chi_{\mathcal{A}}(t) = \chi_{\mathcal{A}'}(t) - \chi_{\mathcal{A}''}(t).$$

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Zaslavsky's theorem ($r(\mathcal{A}) = (-1)^n \chi_{\mathcal{A}}(-1)$) is an immediate consequence of above lemma and the recurrence $r(\mathcal{A}) = r(\mathcal{A}') + r(\mathcal{A}'')$.

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The proof for $b(\mathcal{A})$ is analogous but a little more complicated.

Whitney's theorem

To prove: $\chi_{\mathcal{A}}(t) = \chi_{\mathcal{A}'}(t) - \chi_{\mathcal{A}''}(t)$.

Basic tool (H. Whitney, 1935, for linear arrangements). A subarrangement $\mathcal{B} \subseteq \mathcal{A}$ is **central** if $\bigcap_{H \in \mathcal{B}} H \neq \emptyset$.

Whitney's theorem

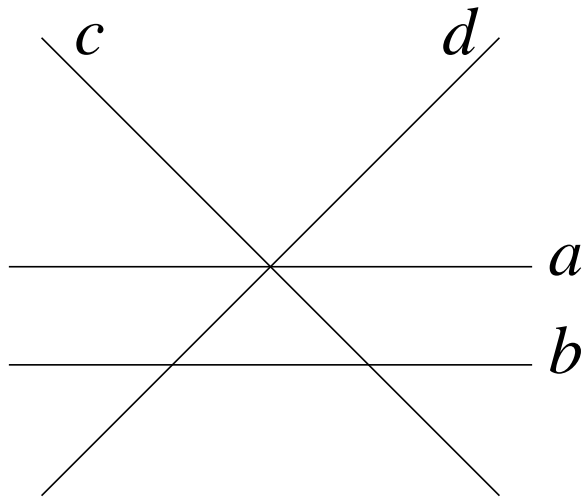
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Theorem. Let \mathcal{A} be an arrangement in an n -dimensional vector space. Then

$$\chi_{\mathcal{A}}(t) = \sum_{\substack{\mathcal{B} \subseteq \mathcal{A} \\ \mathcal{B} \text{ central}}} (-1)^{\#\mathcal{B}} t^{n - \text{rk}(\mathcal{B})}.$$

Example of Whitney's theorem



\mathcal{B}	$\#\mathcal{B}$	$\text{rk}(\mathcal{B})$			
\emptyset	0	0			
a	1	1	bc	2	2
b	1	1	bd	2	2
c	1	1	cd	2	2
d	1	1	acd	3	2
ac	2	2			
ad	2	2			

$$\Rightarrow \chi_{\mathcal{A}}(t) = t^2 - 4t + (5 - 1) = t^2 - 4t + 4.$$

The crosscut theorem

Easy fact: Every interval $[\hat{0}, z]$ of $L(\mathcal{A})$ is a **lattice**, i.e., any two elements x, y have a **meet** (greatest lower bound) $x \wedge y$ and **join** (least upper bound) $x \vee y$.

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Lemma (crosscut theorem for $L(\mathcal{A})$). For all $z \in L(\mathcal{A})$,

$$\mu(z) = \sum_{\substack{\mathcal{B} \subseteq \mathcal{A}_z \\ z = \bigcap_{H \in \mathcal{B}} H}} (-1)^{\#\mathcal{B}}.$$

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Note that $z = \bigcap_{H \in \mathcal{B}} H$ implies that $\text{rk}(\mathcal{B}) = n - \dim z$. Multiply both sides by $t^{\dim(z)}$ and sum over z to obtain

$$\chi_{\mathcal{A}}(t) = \sum_{\substack{\mathcal{B} \subseteq \mathcal{A} \\ \mathcal{B} \text{ central}}} (-1)^{\#\mathcal{B}} t^{n - \text{rk}(\mathcal{B})}. \quad \square$$

Alternative formulation

Later: coefficients of $\chi_{\mathcal{A}}(t)$ **alternate in sign**.
More strongly, if $\text{rk}(x) = i$ then

$$(-1)^i \mu(x) > 0.$$

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Thus:

$$r(\mathcal{A}) = \sum_{x \in L_{\mathcal{A}}} |\mu(x)|$$

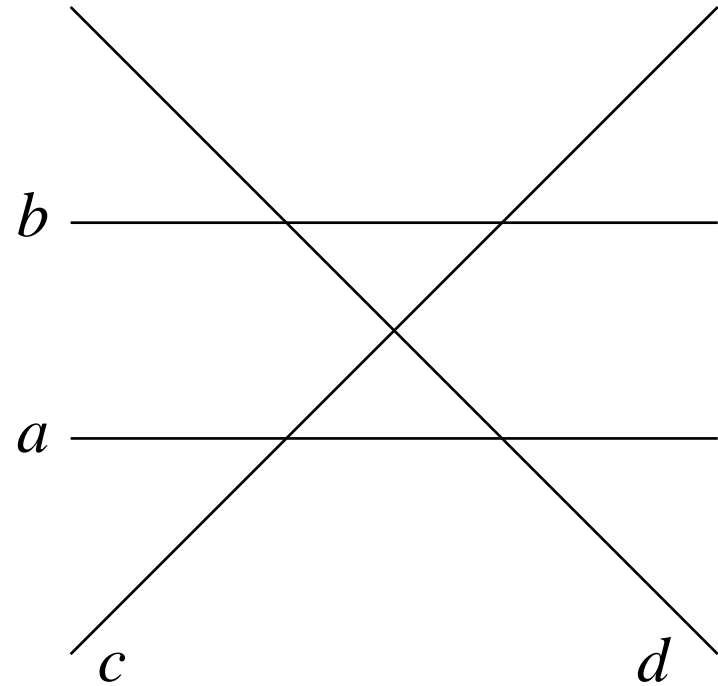
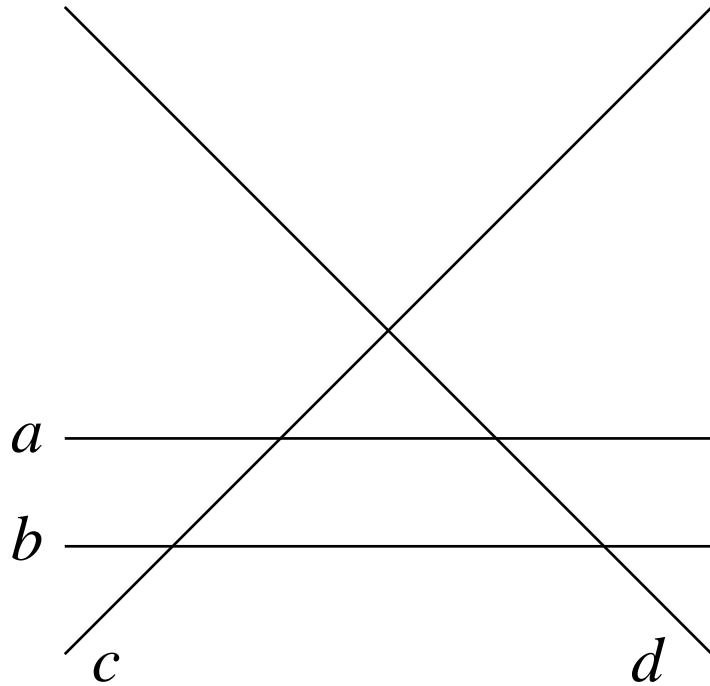
$$b(\mathcal{A}) = \left| \sum_{x \in L_{\mathcal{A}}} \mu(x) \right|.$$

A corollary

Corollary. *Let \mathcal{A} be a real arrangement. Then $r(\mathcal{A})$ and $b(\mathcal{A})$ depend only on $L(\mathcal{A})$.*

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Faces

$\mathcal{R}(\mathcal{A})$: set of regions of \mathcal{A}

Definition. A (closed) **face** of a real arrangement \mathcal{A} is a set

$$\emptyset \neq \mathbf{F} = \overline{R} \cap x,$$

where $R \in \mathcal{R}(\mathcal{A})$, $x \in L(\mathcal{A})$, and \overline{R} = closure of R .

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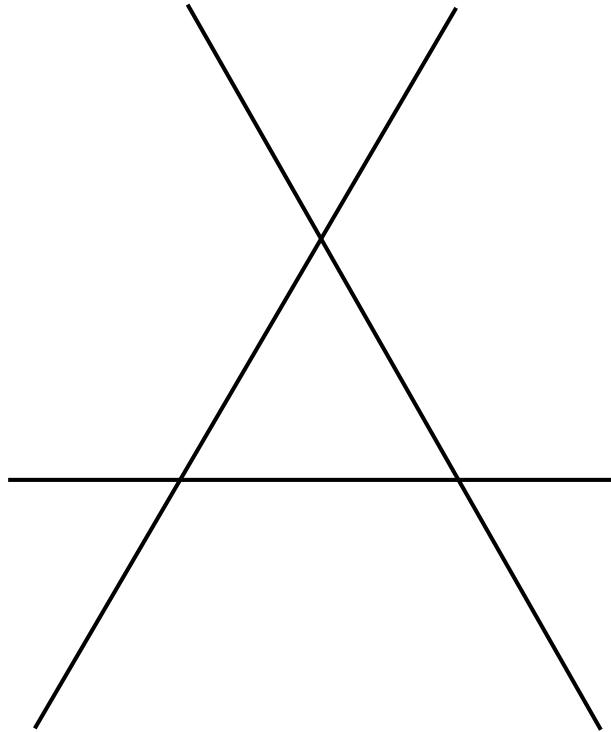
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$f_k(\mathcal{A})$: number of k -dimensional faces (**k -faces**) of \mathcal{A}

Example of $f_i(\mathcal{A})$



$$f_0(\mathcal{A}) = 3, \quad f_1(\mathcal{A}) = 9, \quad f_2(\mathcal{A}) = r(\mathcal{A}) = 7$$

Formula for $f_k(\mathcal{A})$

$$f_k(\mathcal{A}) = \sum_{\substack{x \in L(\mathcal{A}) \\ \text{corank}(x)=k}} \sum_{y \geq x} |\mu(x, y)|$$

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Proof. Easy consequence of Zaslavsky's formula for $r(\mathcal{A})$. \square

Zonotopes

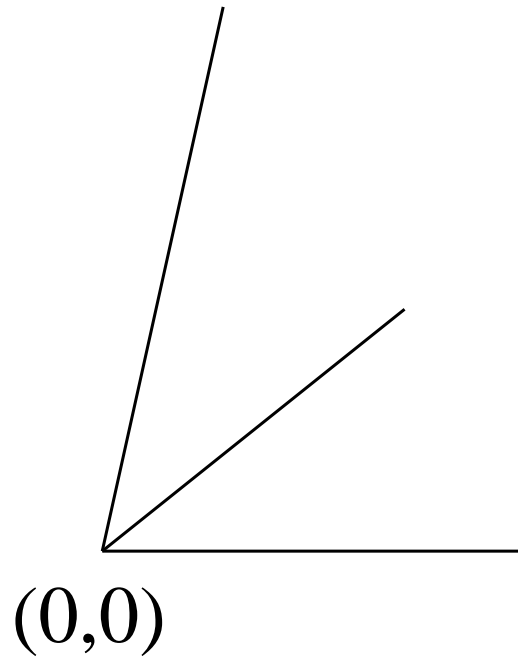
Let $X, Y \subseteq K^n$

Minkowski sum:

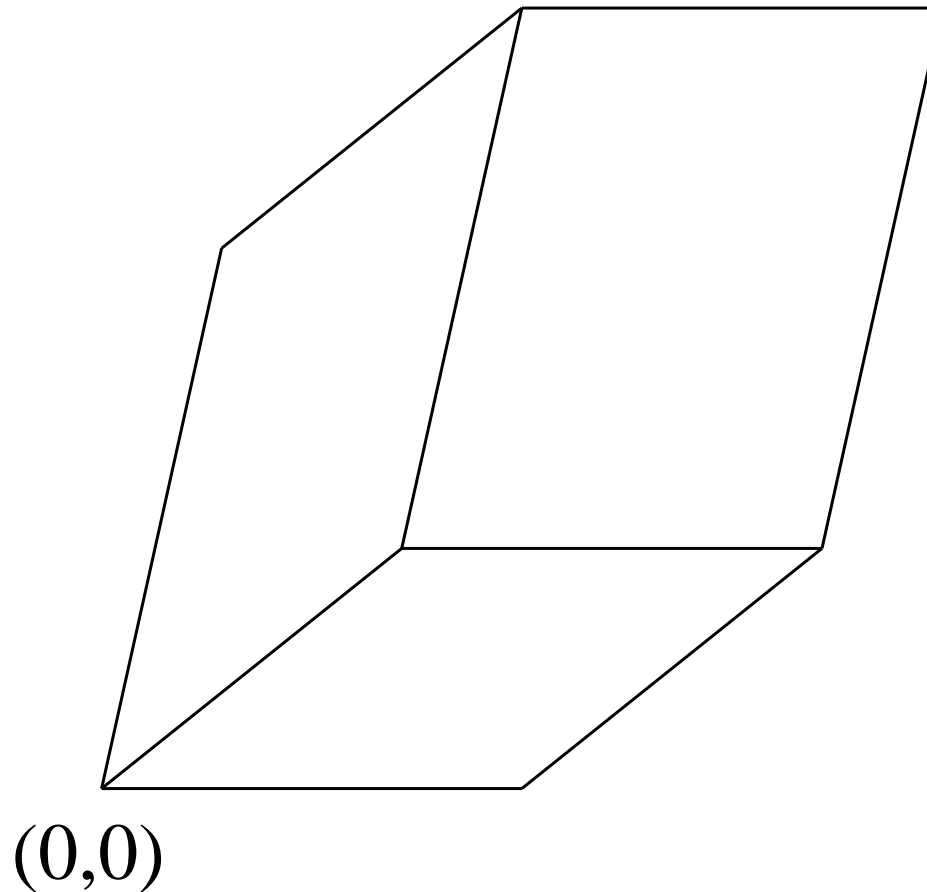
$$\mathbf{X} + \mathbf{Y} = \{x + y : x \in X, y \in Y\}$$

zonotope: a Minkowski sum $L_1 + \cdots + L_k$ of line segments in \mathbb{R}^n

Example of zonotope



Example of zonotope



Characterization of zonotopes

Theorem. *Let \mathcal{P} be a convex polytope. The following are equivalent.*

- \mathcal{P} is a zonotope.
- Every face of \mathcal{P} is centrally-symmetric.
- Every 2-dimensional face of \mathcal{P} is centrally-symmetric.

The zonotope of a real arrangement

\mathcal{A} : a real central arrangement

n_1, \dots, n_k : normals to $H \in \mathcal{A}$

L_i : line segment from 0 to n_i

$Z(\mathcal{A})$: the zonotope $L_1 + \dots + L_k$

Number of faces of $Z(\mathcal{A})$

Theorem. Let $f_i(Z(\mathcal{A}))$ denote the number of i -dimensional faces of $Z(\mathcal{A})$. Then

$$f_i(Z(\mathcal{A})) = f_{n-i}(\mathcal{A}).$$

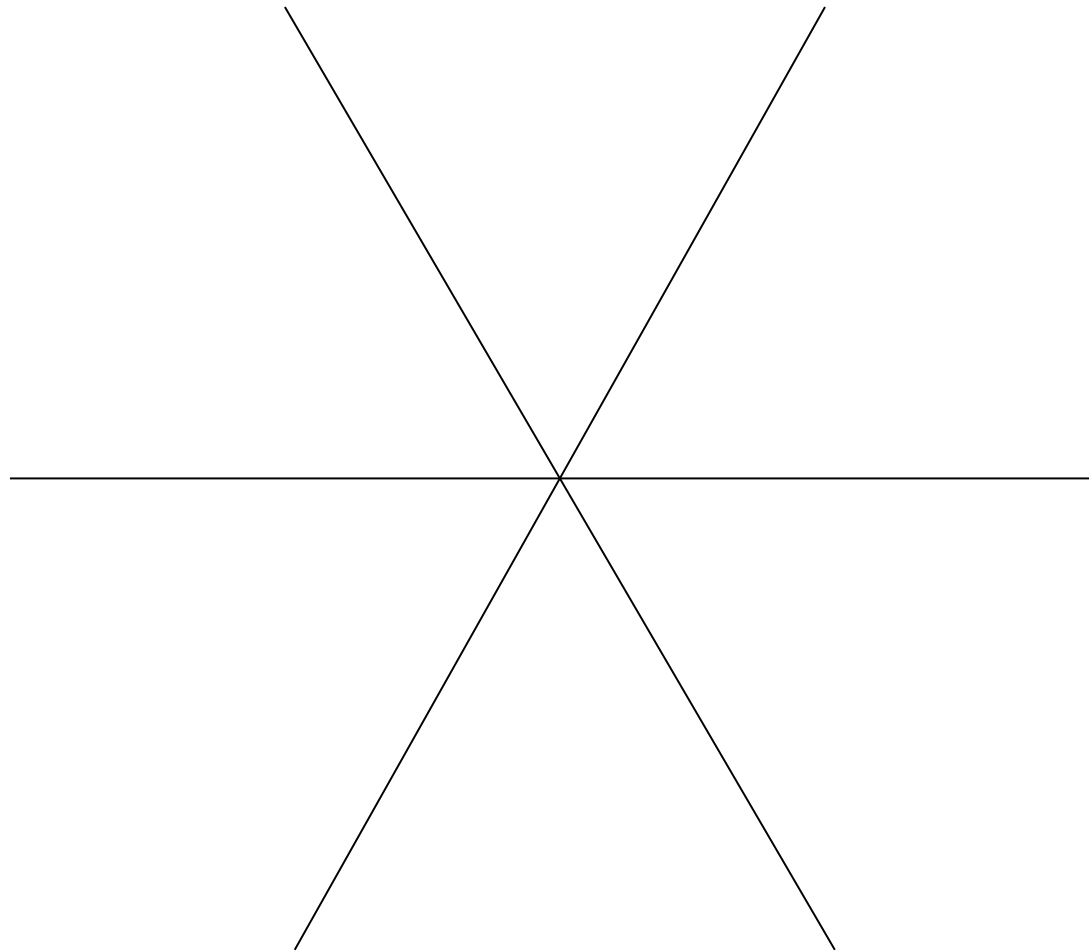
Number of faces of $Z(\mathcal{A})$

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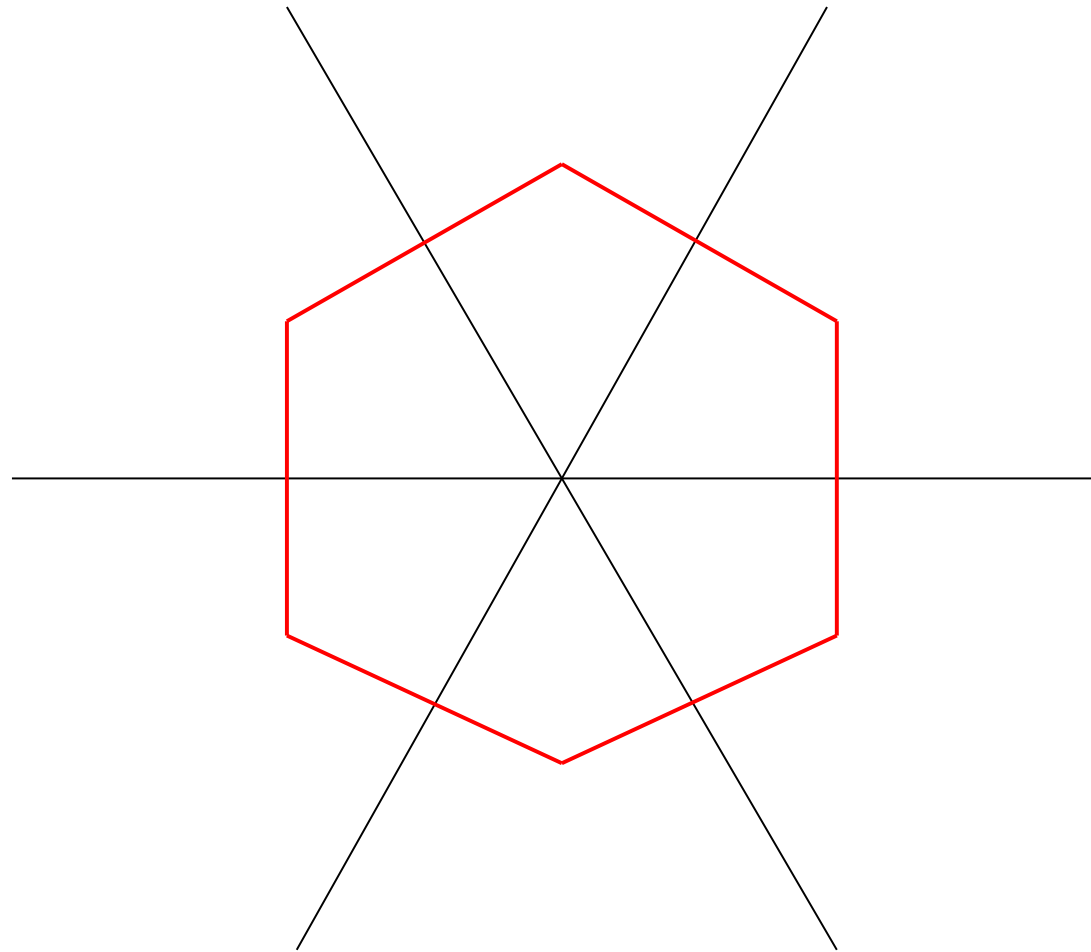
$$f_i(Z(\mathcal{A})) = f_{n-i}(\mathcal{A}).$$

Informally, $Z(\mathcal{A})$ is a “dual object” to \mathcal{A} .

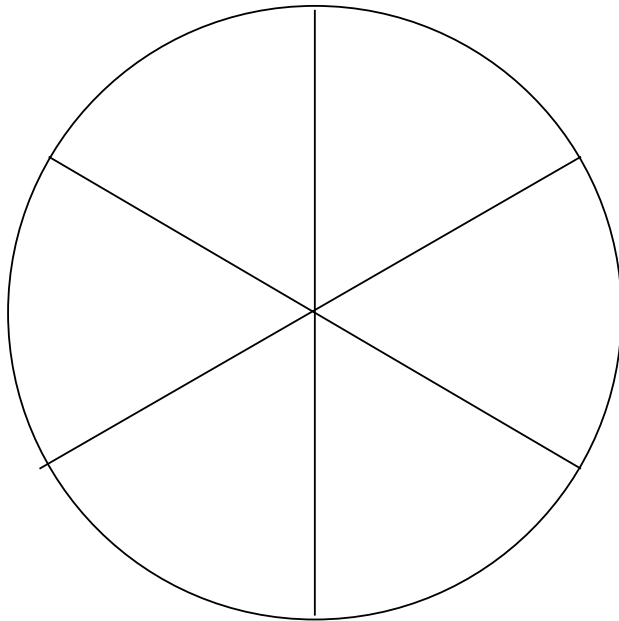
An example of $Z(\mathcal{A})$



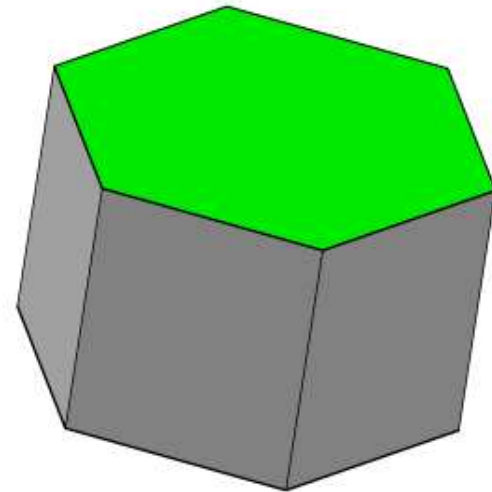
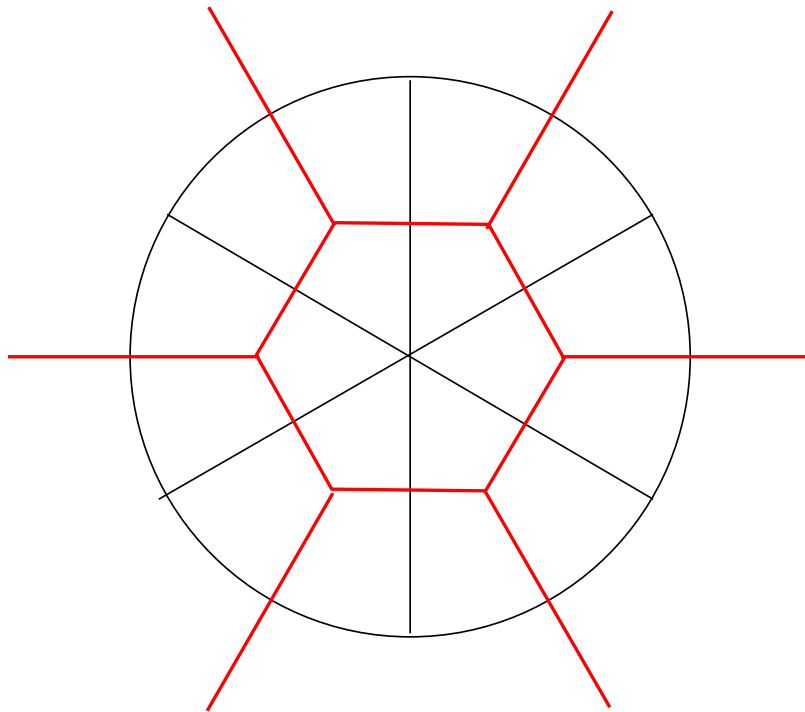
An example of $Z(\mathcal{A})$



Another example

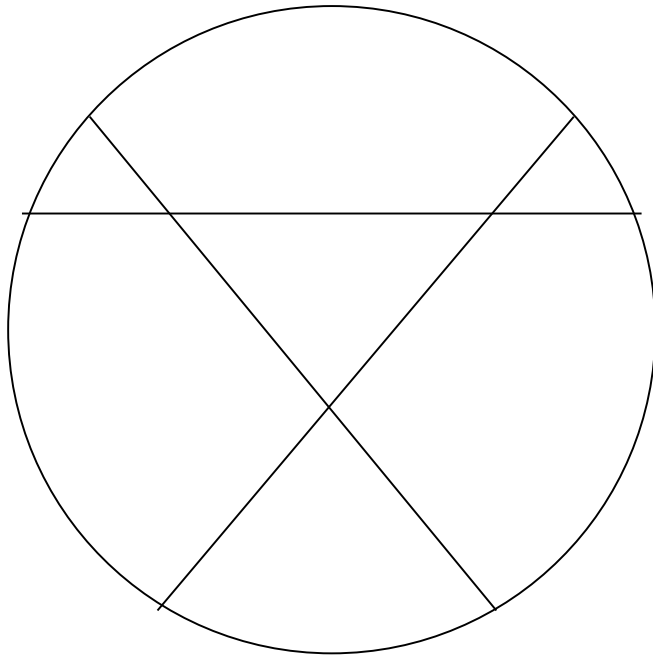


Another example

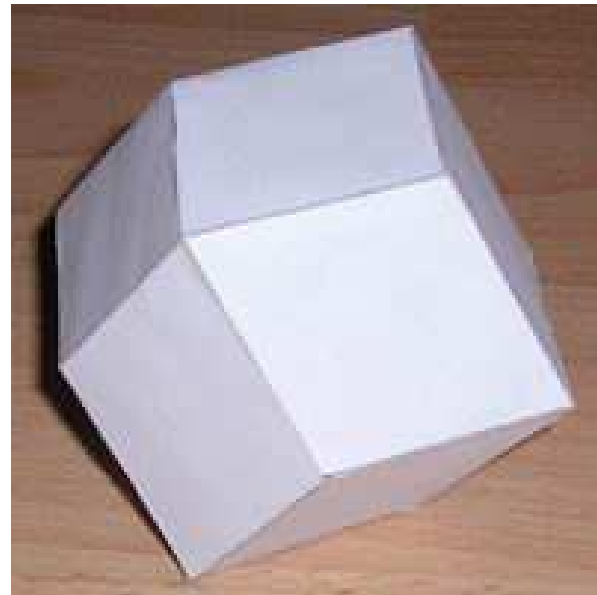
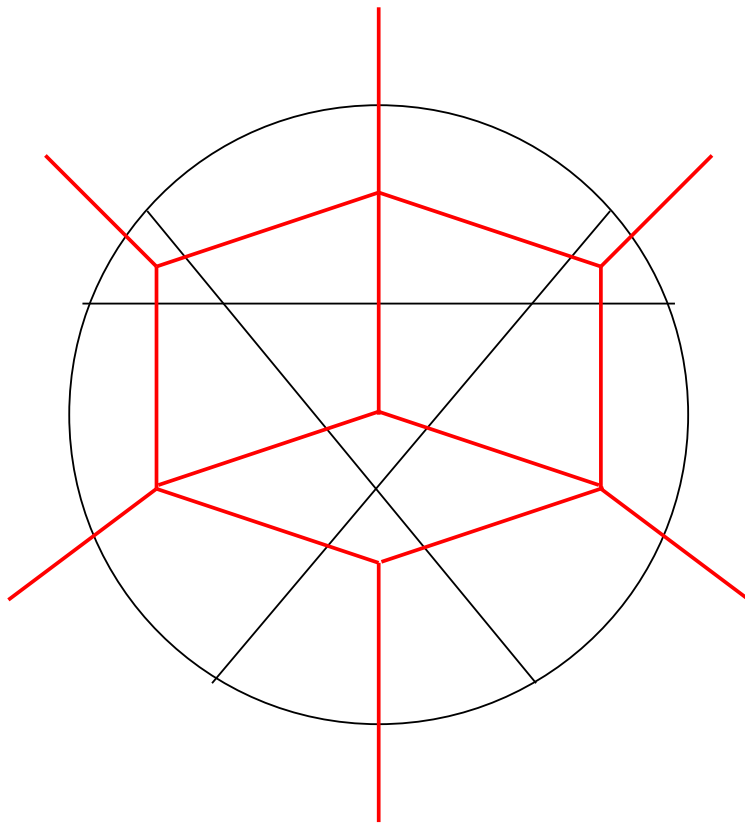


hexagonal prism

Another example



Another example



rhombic dodecahedron

Graphical arrangements

G : graph on vertex set $[n]$ (no loops or multiple edges)

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If $G = K_n$, the **complete graph** on $[n]$, then \mathcal{A}_{K_n} is the **braid arrangement** \mathcal{B}_n .

Set partitions

partition of a finite set S : $\pi = \{B_1, \dots, B_k\}$, such that

$$B_i \neq \emptyset, \quad \bigcup B_i = S, \quad B_i \cap B_j = \emptyset \quad (i \neq j)$$

B_i is a **block** of π .

Π_S : set of partitions of S

Let $\pi, \sigma \in \Pi_S$. Then π is a **refinement** of σ , written $\pi \leq \sigma$, if every block of π is contained in a block of σ .

The bond lattice of G

G : graph on vertex set $[n]$

connected partition of $[n]$: a partition of $[n]$ for which each block induces a connected subgraph of G

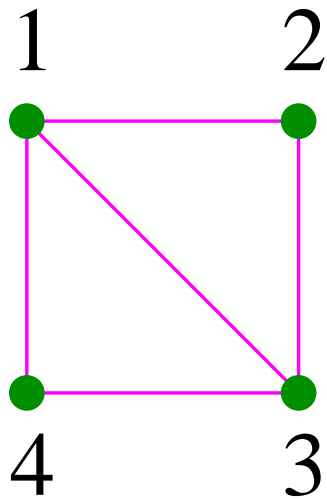
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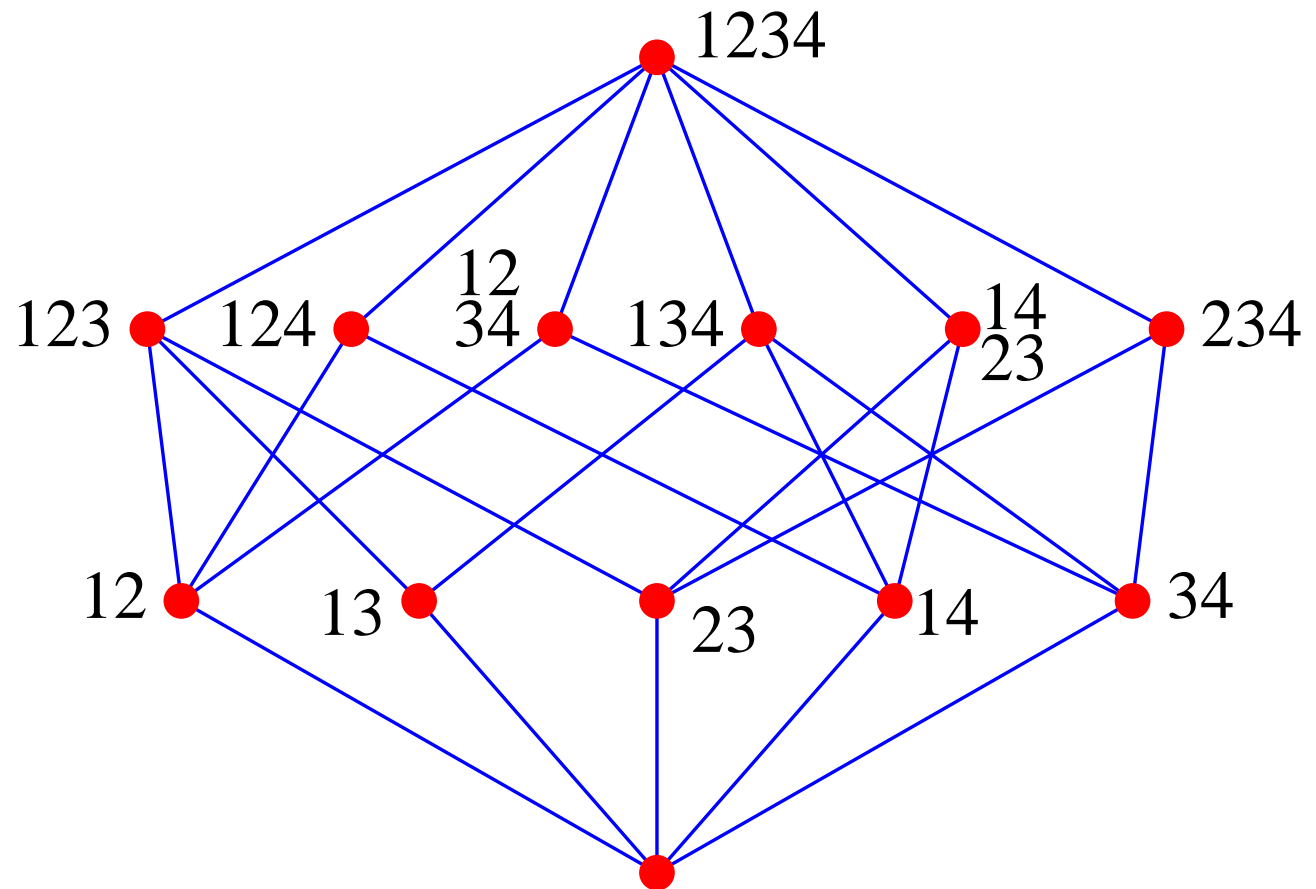
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bond lattice $L(G)$ of G : set of connected partitions of $[n]$, ordered by refinement

Example of bond lattice



G



$L(G)$

Bond lattices and intersection posets

G : graph with bond lattice $L(G)$

\mathcal{A}_G : graphical arrangement

Theorem. $L(G) \cong L(\mathcal{A}(G))$

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Proof. Let H_{ij} be the hyperplane defined by $x_i = x_j$, $ij \in E(G)$. Let $x \in L(\mathcal{A})$. Define vertices $i \sim j$ if $x \subseteq H_{ij}$. Then \sim is an equivalence relation whose equivalence classes form a connected partition of $[n]$, etc. \square

Chromatic polynomial of G

coloring of G is $\kappa: [n] \rightarrow \mathbb{P} = \{1, 2, \dots\}$

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Easy fact: $\chi_G(q) \in \mathbb{Z}[q]$

$$\chi_{\mathcal{A}(G)}(t)$$

Theorem. $\chi_{\mathcal{A}(G)}(t) = \chi_G(t)$

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Proof. Let $\sigma \in L(G)$.

$\chi_{\sigma}(q)$ = number of $f: [n] \rightarrow [q]$ such that:

- a, b in same block of $\sigma \Rightarrow f(a) = f(b)$
- a, b in different blocks, $ab \in E \Rightarrow f(a) \neq f(b)$.

Continuation of proof

Given **any** $f: [n] \rightarrow [q]$, there is a unique $\sigma \in L(G)$ such that f is enumerated by $\chi_\sigma(q)$.
Hence $\forall \pi \in L(G)$,

$$q^{\#\pi} = \sum_{\sigma \geq \pi} \chi_\sigma(q).$$

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Note $\chi_{\hat{0}}(q) = \chi_G(q)$. \square

Characteristic polynomial of \mathcal{B}_n

Recall: $\mathcal{B}_n = \mathcal{A}(K_n)$ (braid arrangement)

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Clearly $\chi_{K_n}(q) = q(q-1) \cdots (q-n+1)$.

$$\Rightarrow \chi_{\mathcal{B}_n}(t) = t(t-1) \cdots (t-n+1).$$

Chordal graphs

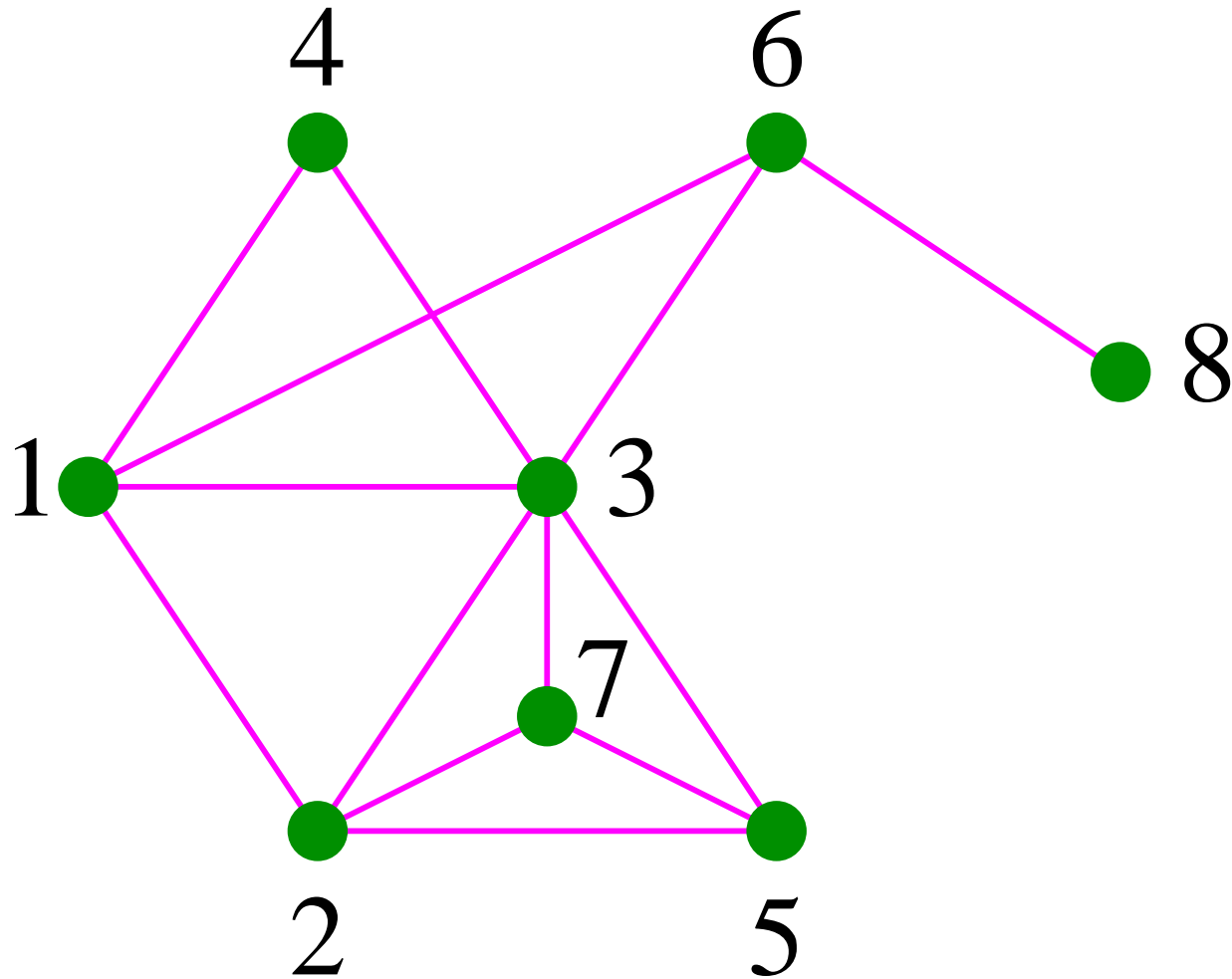
A graph G is **chordal** (**triangulated**, **rigid circuit**) if the vertices can be ordered v_1, \dots, v_n so that for all i , v_i is connected to a **clique** (complete subgraph) of the restriction of G to $\{v_1, \dots, v_{i-1}\}$.

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Known fact: G is chordal if and only if every cycle of length at least four has a chord.

Example of a chordal graph



Chordal graph coloring

Let v_1, \dots, v_n be a vertex ordering so that for all i , v_i is connected to a clique of the restriction G_{i-1} of G to $\{v_1, \dots, v_{i-1}\}$.

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Let a_i be the number of vertices of G_{i-1} to which v_i is connected (so $a_1 = 0$). Once v_1, \dots, v_{i-1} are (properly) colored, there are $q - a_i$ ways to color v_i . Hence

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$$\chi_G(q) = (q - a_1)(q - a_2) \cdots (q - a_n).$$

Acyclic orientations

Orientation of G : assignment of a direction $i \rightarrow j$ or $j \rightarrow i$ to each edge.

Acyclic orientation: an orientation with no directed cycles

$$\chi_G(-1)$$

Given \mathfrak{o} , define

$$R_{\mathfrak{o}} = \{(x_1, \dots, x_n) \in \mathbb{R}^n : x_i < x_j \text{ whenever } i \rightarrow j \text{ in } \mathfrak{o}\}.$$

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ao(G): number of acyclic orientations of G

Theorem. $r(\mathcal{A}_G) = (-1)^n \chi_G(-1) = \text{ao}(G).$

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This proof is due to Greene (1977).

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Goal: interpret $(-1)^i \mu(x, y)$ combinatorially,
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For simplicity we deal only with hyperplane arrangements, though the “right” level of generality is **matroid theory**.

Broken circuits

\mathcal{A} : central arrangement

circuit: a minimal linearly dependent subset of \mathcal{A}

H_1, H_2, \dots, H_m : ordering of \mathcal{A}

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broken circuit complex:

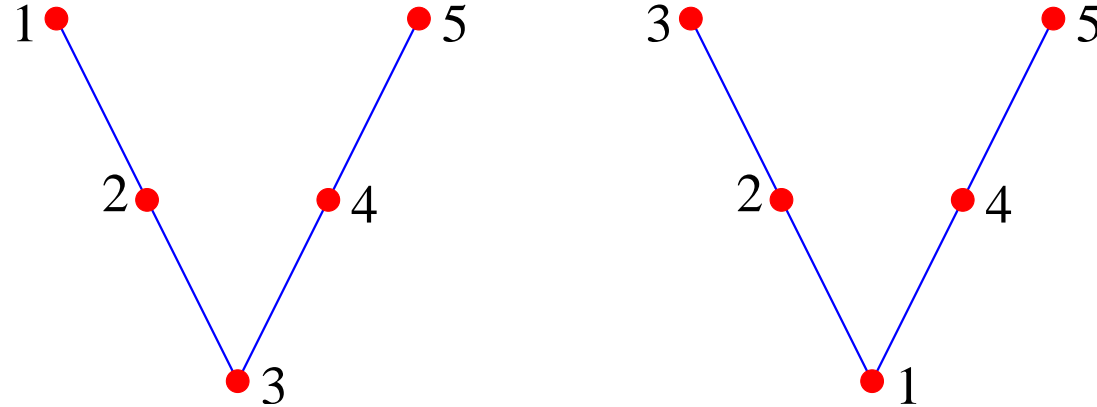
$$\text{BC}(\mathcal{A}) = \{F \subseteq \mathcal{A} : F \text{ contains no broken circuit}\}$$

An example

Note: $\text{BC}(\mathcal{A})$ is a **simplicial complex**, i.e.,
 $F \in \text{BC}(\mathcal{A}), \quad G \subseteq F \Rightarrow G \in \text{BC}(\mathcal{A})$.

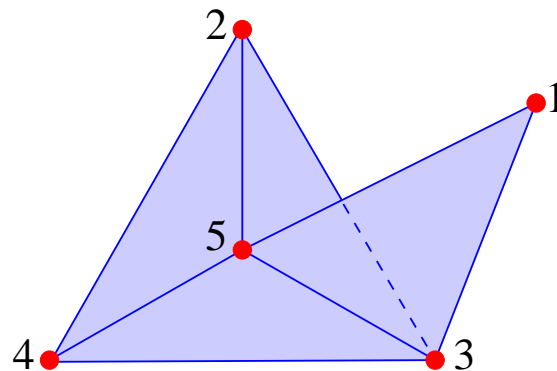
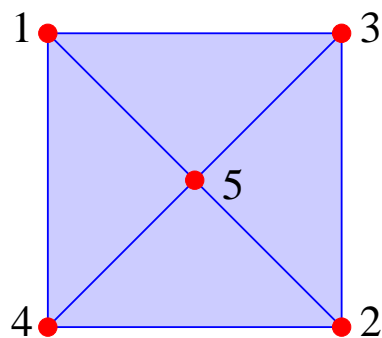
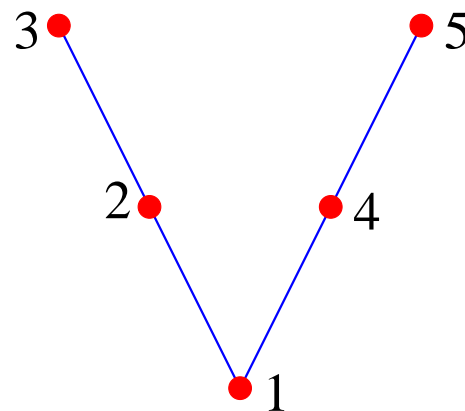
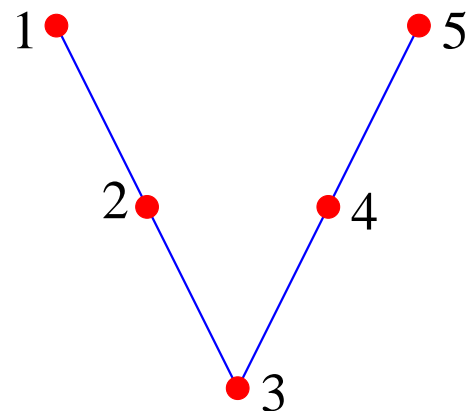
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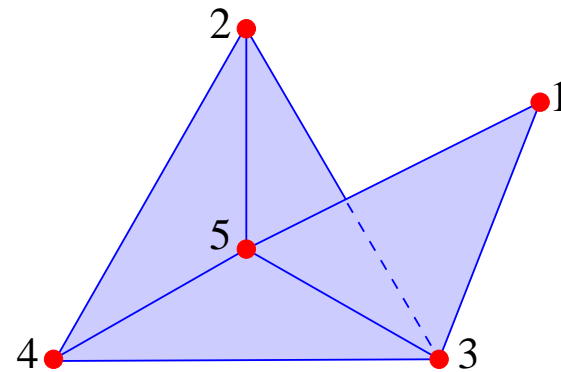
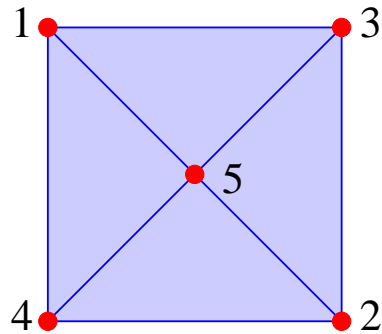


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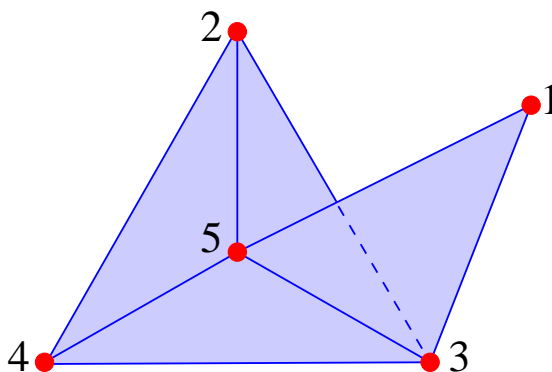
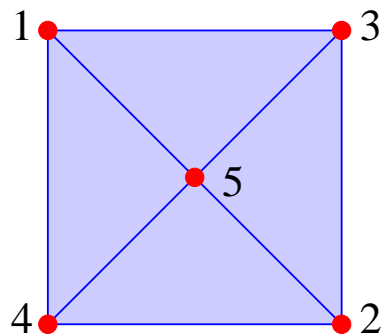
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Example (continued)



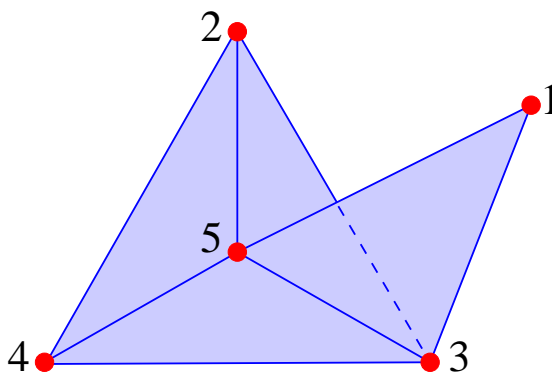
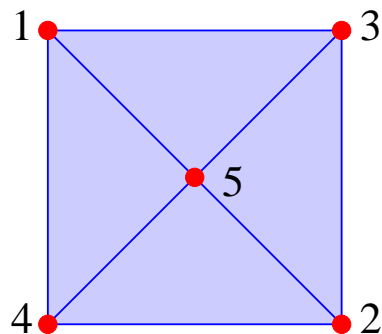
Example (continued)



$f_i = f_i(\text{BC}(\mathcal{A}))$: # i -dim. faces of $\text{BC}(\mathcal{A})$

$$f_{-1} = 1, \quad f_0 = 5, \quad f_1 = 8, \quad f_2 = 4$$

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$$\chi_{\mathcal{A}}(t) = t^3 - 5t^2 + 8t - 4$$

Covers

$$\mathbf{L} = L_{\mathcal{A}}$$

y **covers** x in L : $x < y$, $\nexists x < z < y$

$\mathcal{E}(L)$: edges of Hasse diagram of L , i.e.,

$$\mathcal{E}(L) = \{(x, y) : y \text{ covers } x\}$$

Labelings

$\lambda: \mathcal{E}(L) \rightarrow \mathbb{P}$ is a **labeling** of L

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If $C: x = x_0 < x_1 < \dots < x_k = y$ is a saturated chain from x to y (i.e., each x_{i+1} **covers** x_i), define

$$\lambda(C) = (\lambda(x_0, x_1), \lambda(x_1, x_2), \dots, \lambda(x_{k-1}, x_k))$$

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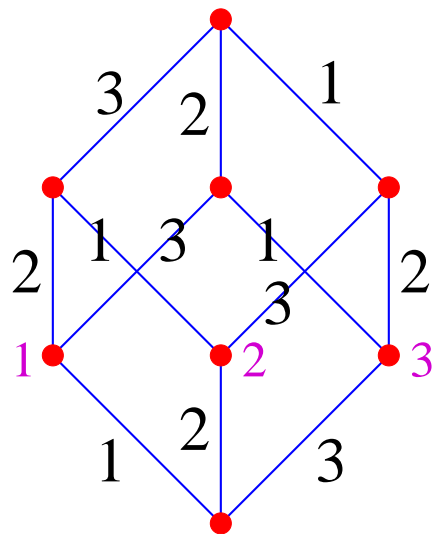
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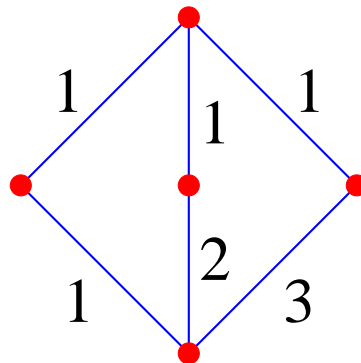
C is **increasing** if

$$\lambda(x_0, x_1) \leq \lambda(x_1, x_2) \leq \dots \leq \lambda(x_{k-1}, x_k).$$

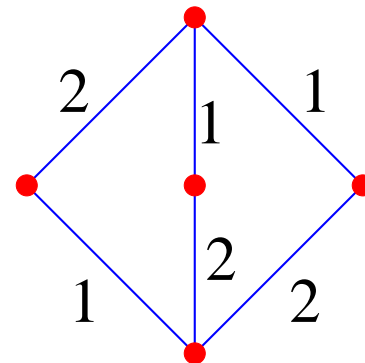
E-labelings



(a)

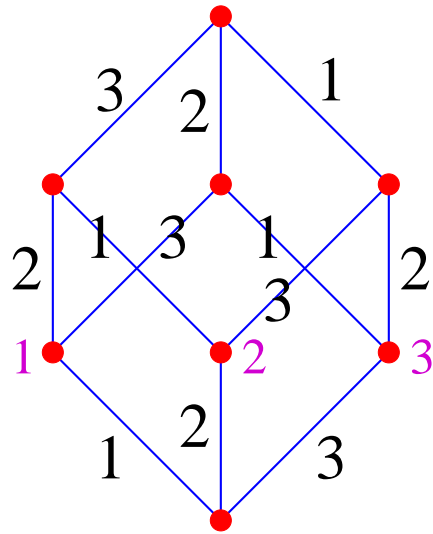


(b)

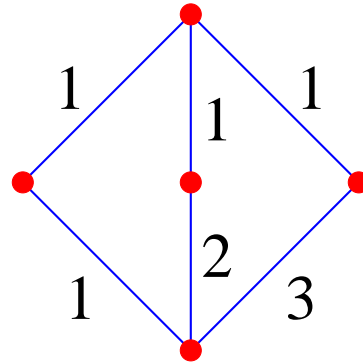


(c)

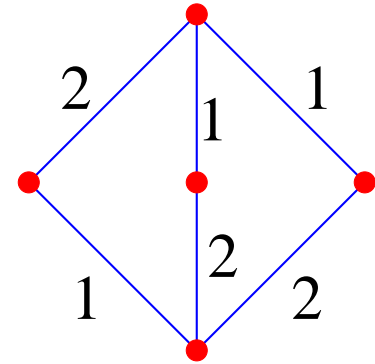
E-labelings



(a)



(b)



(c)

E-labeling: a labeling for which every interval $[x, y]$ has a **unique** increasing chain.

Labeling and Möbius functions

Theorem. Let λ be an E -labeling of L , and let $x \leq y$ in L , $\text{rank}(x, y) = k$. Then $(-1)^k \mu(x, y)$ is equal to the number of strictly decreasing saturated chains from x to y , i.e.,

$$(-1)^k \mu(x, y) = \#\{x = x_0 < x_1 < \cdots < x_k = y : \\ \lambda(x_0, x_1) > \lambda(x_1, x_2) > \cdots > \lambda(x_{k-1}, x_k)\}.$$

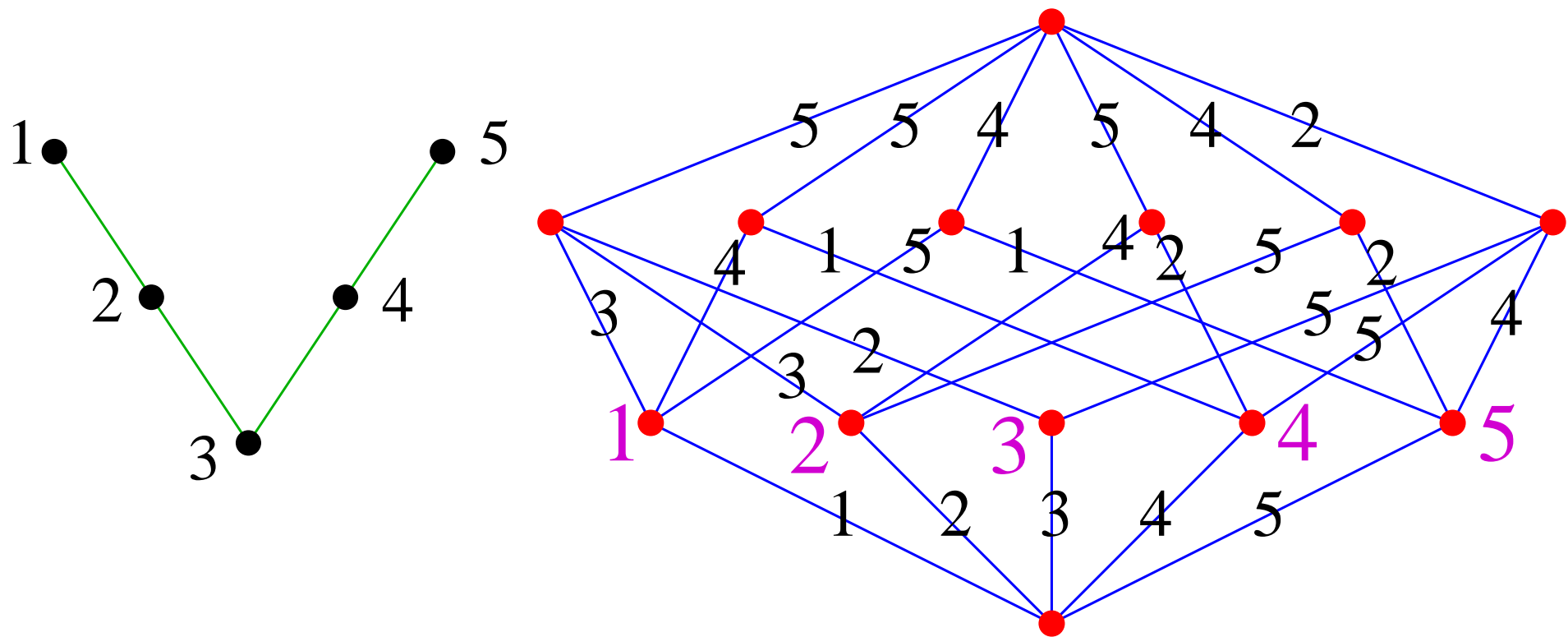
Labeling $L(\mathcal{A})$

H_1, \dots, H_m : ordering of \mathcal{A} (as before)

If y covers x in $L(\mathcal{A})$ then define

$$\tilde{\lambda}(x, y) = \max\{i : x \vee H_i = y\}.$$

Example of λ



Properties of λ

Claim 1. Define $\lambda: \mathcal{E}(L(\mathcal{A})) \rightarrow \mathbb{P}$ by

$$\lambda(x, y) = m + 1 - \tilde{\lambda}(x, y).$$

Then λ is an E -labeling.

Properties of λ

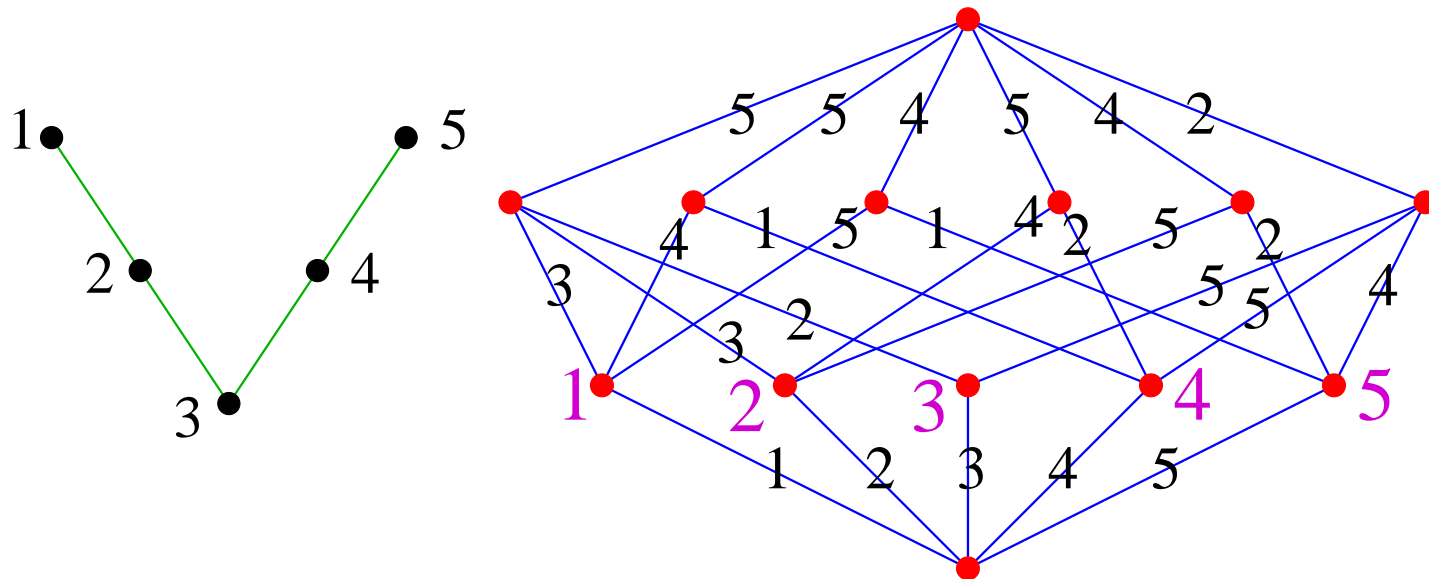
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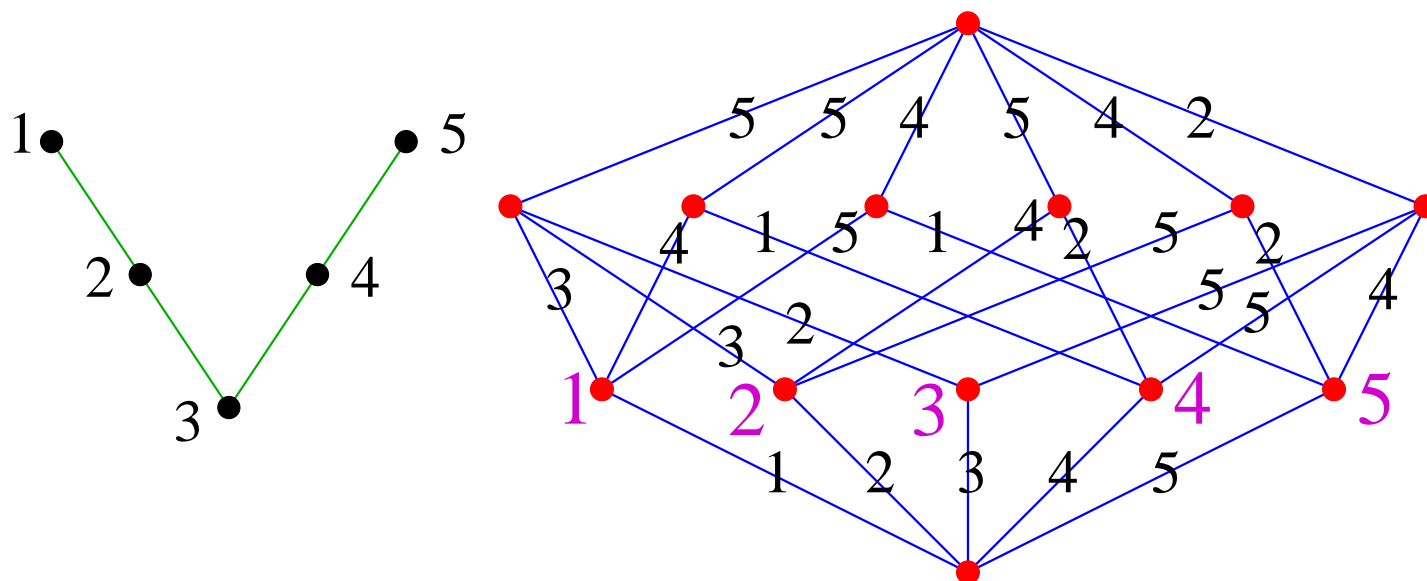
Then λ is an E -labeling.

Claim 2. The broken circuit complex $BC(M)$ consists of all chain labels $\lambda(C)$ (regarded as a **set**), where C is an increasing saturated chain from $\hat{0}$ to some $x \in L(M)$. Moreover, all such $\lambda(C)$ are distinct.

Example of Claim 2.



Example of Claim 2.



broken circuits : 12, 34, 124

$$\text{BC}(\mathcal{A}) = \{\emptyset, 1, 2, 3, 4, 5, 13, 14, 15, 23, 24, 25, 35, 45, 135, 145, 235, 245\}$$

Broken circuit theorem

Immediate consequence of Claims 1 and 2:

Theorem. $\chi_{\mathcal{A}}(t) = \sum_{F \in \text{BC}(\mathcal{A})} (-1)^{\#F} t^{n - \#F}$

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Theorem. $\chi_{\mathcal{A}}(t) = \sum_{F \in \text{BC}(\mathcal{A})} (-1)^{\#F} t^{n-\#F}$

Corollary. *The coefficients of $\chi_{\mathcal{A}}(t)$ **alternate in sign**, i.e., $\chi_{\mathcal{A}}(t) = t^n - a_1 t^{n-1} + a_2 t^{n-2} - \dots$, where $a_i \geq 0$. In fact*

$$(-1)^i \mu(x, y) > 0, \quad \text{where } i = \text{rank}(x, y).$$

A glimpse of topology

$[x, y]$: (finite) interval in a poset P

c_i : number of chains $x = x_0 < x_1 < \cdots < x_i = y$

Note. $c_0 = 0$ unless $x = y$.

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Philip Hall's theorem (1936).

$$\mu(x, y) = c_0 - c_1 + c_2 - \cdots$$

The order complex

P : a poset

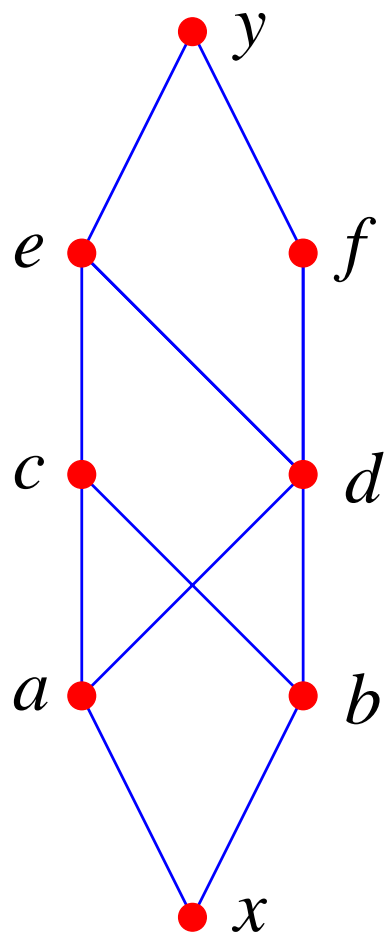
order complex of P :

$$\Delta(P) = \{\text{chains of } P\},$$

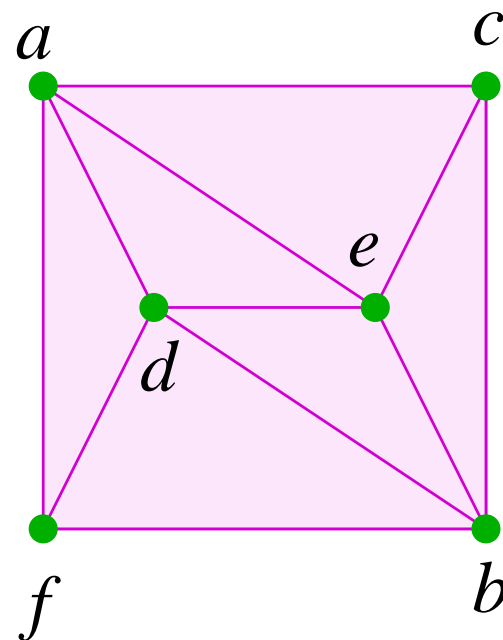
an abstract simplicial complex.

Write $\Delta(x, y)$ for the order complex of the **open** interval $(x, y) = \{z \in P : x < z < y\}$.

Example of an order complex



P



$\Delta(P)$

Euler characteristic

Δ : finite simplicial complex

$f_i = \#$ i -dimensional faces of Δ

Note: $f_{-1} = 1$ unless $\Delta = \emptyset$.

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reduced Euler characteristic:

$\tilde{\chi}(\Delta) = -f_{-1} + f_0 - f_1 + f_2 - \cdots$

Note: $\tilde{\chi}(\Delta) = \chi(\Delta) - 1$ unless $\Delta = \emptyset$.

Philip Hall's theorem restated

Theorem. *For $x < y$ in a finite poset,*

$$\mu(x, y) = \tilde{\chi}(\Delta(x, y)).$$

Philip Hall's theorem restated

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Recall for any finite simplicial complex Δ ,

$$\tilde{\chi}(\Delta) = \sum_j (-1)^j \dim \tilde{H}_j(\Delta; K),$$

where $\tilde{H}_j(\Delta; K)$ denotes reduced simplicial homology over the field K .

A topological question

For $x < y$ in $L(\mathcal{A})$, with $i = \text{rank}(x, y)$, we have

$$\textcolor{red}{d} := \dim \Delta(x, y) = i - 2.$$

In particular, $(-1)^d = (-1)^i$.

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Is there a topological reason for this?

Folkman's theorem

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Theorem (Folkman, 1966).

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Early result in **topological combinatorics**.

Cohen-Macaulay posets

A finite poset P is **Cohen-Macaulay** (over K) if after adjoining a top and bottom element to P , every interval $[x, y]$ satisfies:

If $d = \dim \Delta(x, y)$ then $\tilde{H}_j(\Delta(x, y); k) = 0$ whenever $j \neq d$.

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Folkman's theorem, restated. *If \mathcal{A} is central then $L(\mathcal{A})$ is Cohen-Macaulay.*

Modular elements

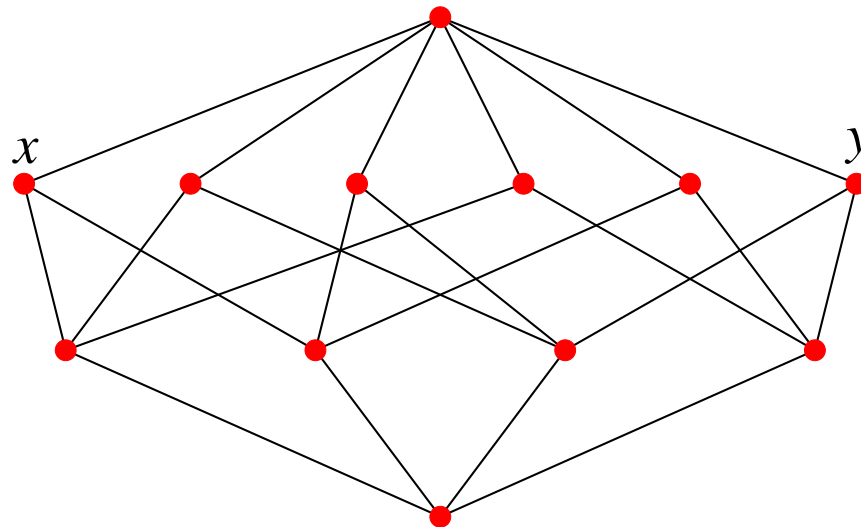
Let \mathcal{A} be central. An element $x \in L(\mathcal{A})$ is **modular** if for all $y \in L$ we have

$$\text{rk}(x) + \text{rk}(y) = \text{rk}(x \wedge y) + \text{rk}(x \vee y).$$

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x is **not** modular: $\text{rk}(x) + \text{rk}(y) = 2 + 2 = 4$,
 $\text{rk}(x \wedge y) + \text{rk}(x \vee y) = 0 + 3 = 3$

Simple properties

Easy: $\hat{0} = K^n$, $\hat{1} = \bigcap_{H \in \mathcal{A}} H$ (the top element),
and each $H \in \mathcal{A}$ is modular.

More properties

$x, y \in L(\mathcal{A})$ are **complements** if $x \wedge y = \hat{0}$,
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Theorem. *Let $r = \text{rk}(\mathcal{A})$. Let $x \in L$. The following four conditions are equivalent.*

- (i) *x is a modular element of L .*
- (ii) *If $x \wedge y = \hat{0}$, then $\text{rk}(x) + \text{rk}(y) = \text{rk}(x \vee y)$.*
- (iii) *If x and y are complements, then $\text{rk}(x) + \text{rk}(y) = n$.*
- (iv) *All complements of x are incomparable.*

Two additional results

Theorem.

- (a) (transitivity of modularity) *If x is a modular element of L and y is modular in the interval $[\hat{0}, x]$, then y is a modular element of L .*
- (b) *If x and y are modular elements of L , then $x \wedge y$ is also modular.*

Modular element factorization thm.

Theorem. *Let z be a modular element of $L(\mathcal{A})$, \mathcal{A} central of rank r . Write $\chi_z(t) = \chi_{[\hat{0}, z]}(t)$. Then*

$$\chi_L(t) = \chi_z(t) \left[\sum_{y: y \wedge z = \hat{0}} \mu_L(y) t^{r - \text{rk}(y) - \text{rk}(z)} \right].$$

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Since each $H \in \mathcal{A}$ is modular in $L(\mathcal{A})$, we get:

Corollary. For all $H \in \mathcal{A}$,

$$\chi_L(t) = (t - 1) \sum_{y \wedge H = \hat{0}} \mu(y) t^{n-1 - \text{rk}(y)}.$$

Supersolvability

A central arrangement \mathcal{A} (or $L(\mathcal{A})$) is **supersolvable** if $L(\mathcal{A})$ has a maximal chain $\hat{0} = x_0 < x_1 < \cdots < x_r = \hat{1}$ of modular elements x_i .

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In this case, let

$$a_i = \#\{H \in \mathcal{A} : H \leq x_i, H \not\leq x_{i-1}\}.$$

Corollary. *If \mathcal{A} is supersolvable, then*

$$\chi_{\mathcal{A}}(t) = t^{n-r} (t - a_1)(t - a_2) \cdots (t - a_n).$$

Chordal graphs, revisited

For what graphs G is \mathcal{A}_G supersolvable?

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If v_i is connected to a_i vertices of G_{i-1} , then

$$\chi_G(q) = (q - a_1)(q - a_2) \cdots (q - a_n).$$

Supersolvable graphs

Suggests that

G chordal $\Rightarrow G$ (or \mathcal{A}_G) supersolvable.

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In fact:

Theorem. *G is chordal if and only if \mathcal{A}_G is supersolvable.*

Free arrangements

Terao (1980) defined **free arrangements** \mathcal{A} and proved

$$\chi_{\mathcal{A}}(t) = (t - a_1) \cdots (t - a_n),$$

where $a_i \in \{0, 1, 2, \dots\}$. (Definition not given here.)

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Supersolvable arrangements are free.

Open: is freeness of \mathcal{A} a combinatorial property? That is, does it just depend on $\chi_{\mathcal{A}}(t)$?

Finite fields and good reduction

\mathcal{A} : arrangement over \mathbb{Q}

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\mathcal{A}_q has **good reduction** if $L_{\mathcal{A}} \cong L_{\mathcal{A}_q}$.

Almost always good reduction

Example. $\mathcal{A} = \{2, 10\}$: affine arrangement in $\mathbb{Q}^1 = \mathbb{Q}$. Good reduction $\Leftrightarrow p \neq 2, 5$.

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Proof idea. Consider minors of the coefficient matrix, etc. \square

The finite field method

Theorem. *Let \mathcal{A} be an arrangement in \mathbb{Q}^n , and suppose that $L(\mathcal{A}) \cong L(\mathcal{A}_q)$ for some prime power q . Then*

$$\begin{aligned}\chi_{\mathcal{A}}(q) &= \# \left(\mathbb{F}_q^n - \bigcup_{H \in \mathcal{A}_q} H \right) \\ &= q^n - \# \bigcup_{H \in \mathcal{A}_q} H.\end{aligned}$$

Proof

Let $x \in L(\mathcal{A}_q)$ so $\#x = q^{\dim(x)}$ (computed either over \mathbb{Q} or F_q). Define $f, g : L(\mathcal{A}_q) \rightarrow \mathbb{Z}$ by

$$f(x) = \#x$$

$$g(x) = \# \left(x - \bigcup_{y>x} y \right)$$

$$\Rightarrow g(\hat{0}) = g(\mathbb{F}_q^n) = \# \left(\mathbb{F}_q^n - \bigcup_{H \in \mathcal{A}_q} H \right).$$

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$$x = \hat{0} \Rightarrow g(\hat{0}) = \sum_y \mu(y) q^{\dim(y)} = \chi_{\mathcal{A}}(q)$$

Graphical arrangements

G : graph on vertex set $1, 2, \dots, n$

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The braid arrangement $\mathcal{B}(B_n)$

$$x_i - x_j = 0, \quad 1 \leq i < j \leq n$$

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Choose α_1 in $q - 1$ ways, then α_2 in $q - 3$ ways, etc.

Characteristic polynomial of $\mathcal{B}(B_n)$

$$\Rightarrow \chi_{\mathcal{B}(B_n)}(q) = (q - 1)(q - 3) \cdots (q - 2n + 1)$$

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In fact, $\mathcal{B}(B_n)$ is supersolvable.

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Exercise: If $n \geq 3$ then

$$\chi_{\mathcal{B}(D_n)} = (q-1)(q-3) \cdots (q-2n+3) \cdot (q-n+1).$$

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Exercise: If $n \geq 3$ then

$$\chi_{\mathcal{B}(D_n)} = (q-1)(q-3) \cdots (q-2n+3) \cdot (q-n+1).$$

Not supersolvable ($n \geq 4$), but it is free.

The Shi arrangement

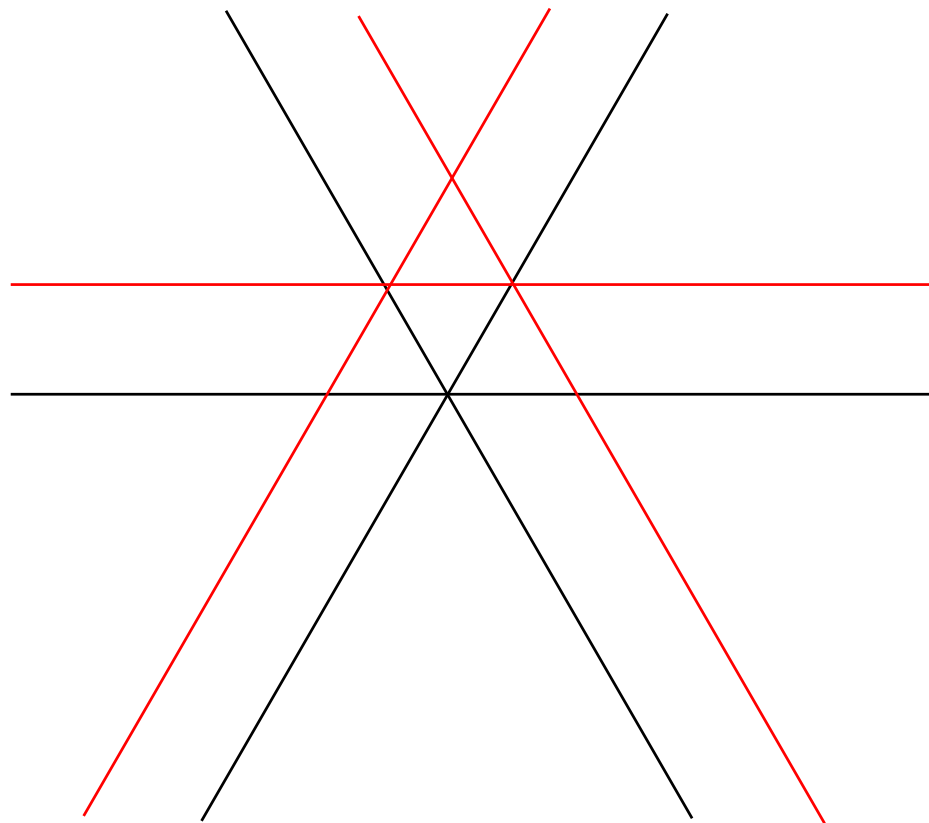
$$\mathcal{S}_n : x_i - x_j = 0, 1, \quad 1 \leq i < j \leq n$$

$$\dim \mathcal{S}_n = n, \quad \text{rk } \mathcal{S}_n = n - 1, \quad \#\mathcal{S}_n = n(n - 1)$$

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Characteristic polynomial of \mathcal{S}_n

Theorem. $\chi_{\mathcal{S}_n}(t) = t(t - n)^{n-1}$, so

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Proof. Finite field method \Rightarrow

$$\chi_{\mathcal{S}_n}(p) = \#\{(\alpha_1, \dots, \alpha_n) \in \mathbb{F}_p^n :$$

$$i < j \Rightarrow \alpha_i \neq \alpha_j \text{ and } \alpha_i \neq \alpha_j + 1\},$$

for $p \gg 0$ (actually, all p).

Proof continued

Choose $\pi = (B_1, \dots, B_{p-n})$ such that

$$\bigcup B_i = [n], \quad B_i \cap B_j = \emptyset \text{ if } i \neq j, \quad 1 \in B_1.$$

Proof continued

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For $2 \leq k \leq n$ there are $p - n$ choices for i such that $k \in B_i$, so $(p - n)^{n-1}$ choices in all.

Circular placement of \mathbb{F}_p

Arrange the elements of \mathbb{F}_p clockwise on a circle.

Place $1, 2, \dots, n$ on some n of these points as follows.

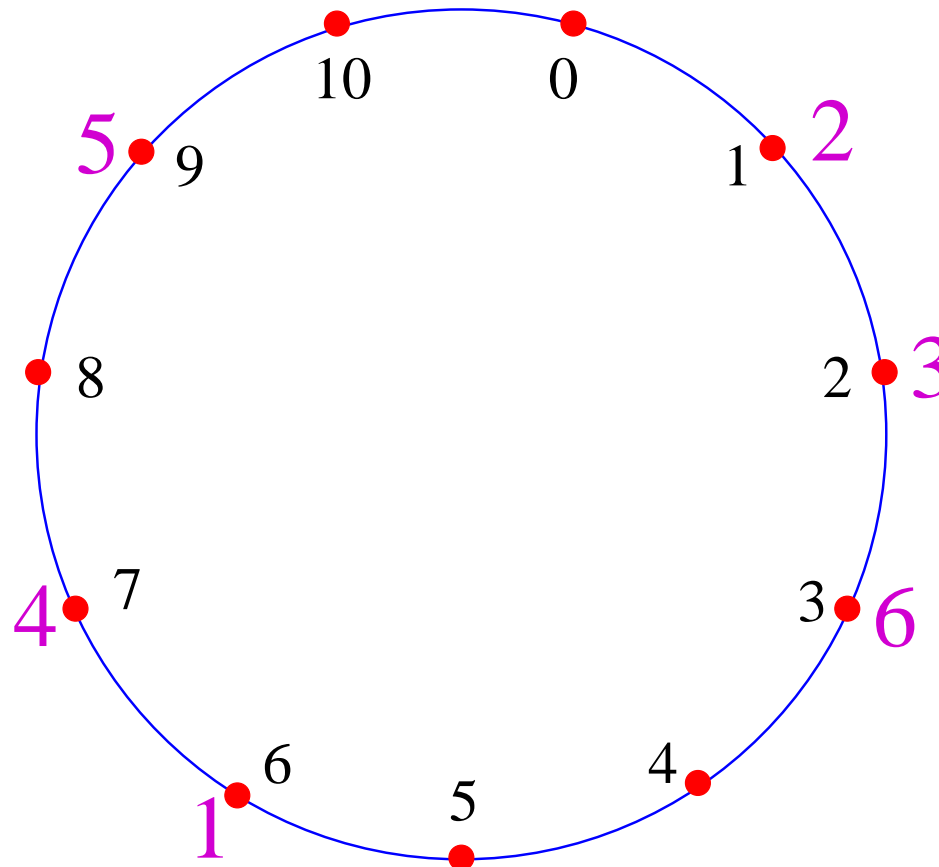
Place elements of B_1 consecutively (clockwise) in increasing order with 1 placed at some element $\alpha_1 \in \mathbb{F}_p$.

Skip a space and place the elements of B_2 consecutively in increasing order.

Skip another space and place the elements of B_3 consecutively in increasing order, etc.

Example for $p = 11, n = 6$

$$\pi = (\{1, 4\}, \{5\}, \emptyset, \{2, 3, 6\}, \emptyset)$$



Conclusion of proof

α_i : position (element of \mathbb{F}_p) at which i was placed

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$$(\alpha_1, \dots, \alpha_6) = (6, 1, 2, 7, 9, 3) \in \mathbb{F}_{11}^6$$

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Gives bijection

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$(p - n)^{n-1}$ choices for π and p choices for α_1 , so

$$\chi_{\mathcal{S}_n}(p) = p(p - n)^{n-1}.$$

The Catalan arrangement

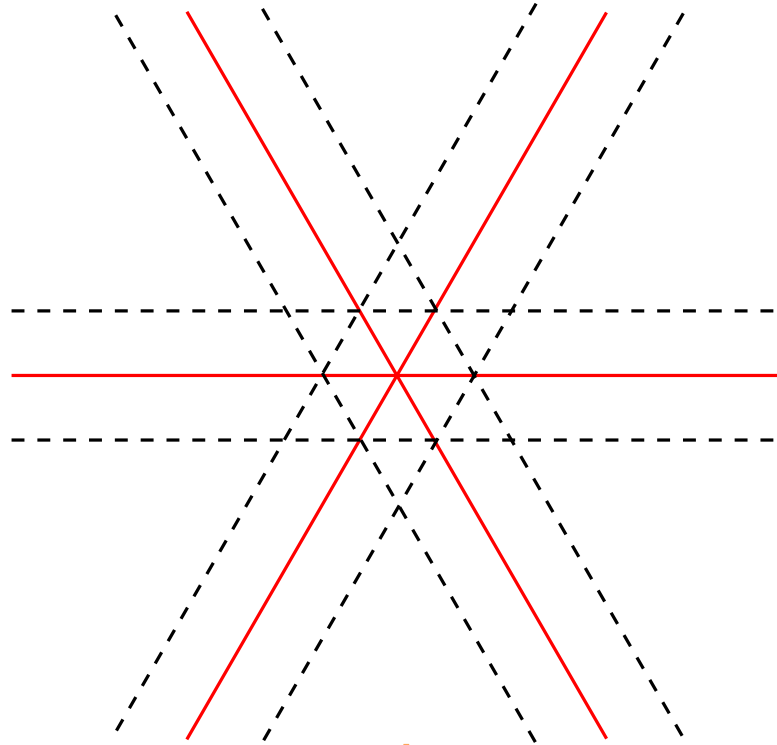
$$\mathcal{C}_n : x_i - x_j = 0, -1, 1, \quad 1 \leq i < j \leq n$$

$$\dim \mathcal{C}_n = n, \quad \text{rk } \mathcal{C}_n = n - 1, \quad \#\mathcal{S}_n = 3 \binom{n}{2}$$

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Char. poly. of Catalan arrangement

Theorem.

$$\chi_{\mathcal{C}_n}(t) = t(t-n-1)(t-n-2)(t-n-3) \cdots (t-2n+1),$$

so

$$r(\mathcal{C}_n) = n!C_n, \quad b(\mathcal{C}_n) = n!C_{n-1},$$

where $C_m = \frac{1}{m+1} \binom{2m}{m}$ (**Catalan number**).

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Each region of the braid arrangement \mathcal{B}_n contains C_n regions and C_{n-1} relatively bounded regions of the Catalan arrangement \mathcal{C}_n .

Catalan numbers

≥ 172 combinatorial interpretations of C_n at

`math.mit.edu/~rstan/ec`

The Linial arrangement

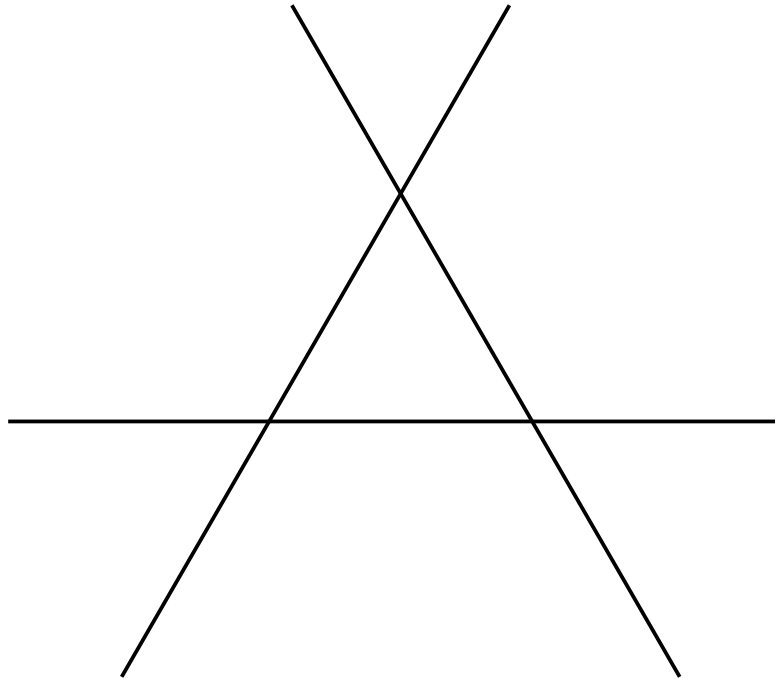
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Char. poly. of Linial arrangement

Theorem. $\chi_{\mathcal{L}_n}(t) = \frac{t}{2^n} \sum_{k=1}^n \binom{n}{k} (t - k)^{n-1},$

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$$r(\mathcal{L}_n) = \frac{1}{2^n} \sum_{k=1}^n \binom{n}{k} (k + 1)^{n-1}$$

$$b(\mathcal{L}_n) = \frac{1}{2^n} \sum_{k=1}^n \binom{n}{k} (k - 1)^{n-1}$$

Two proofs

Postnikov: (difficult) proof using Whitney's theorem

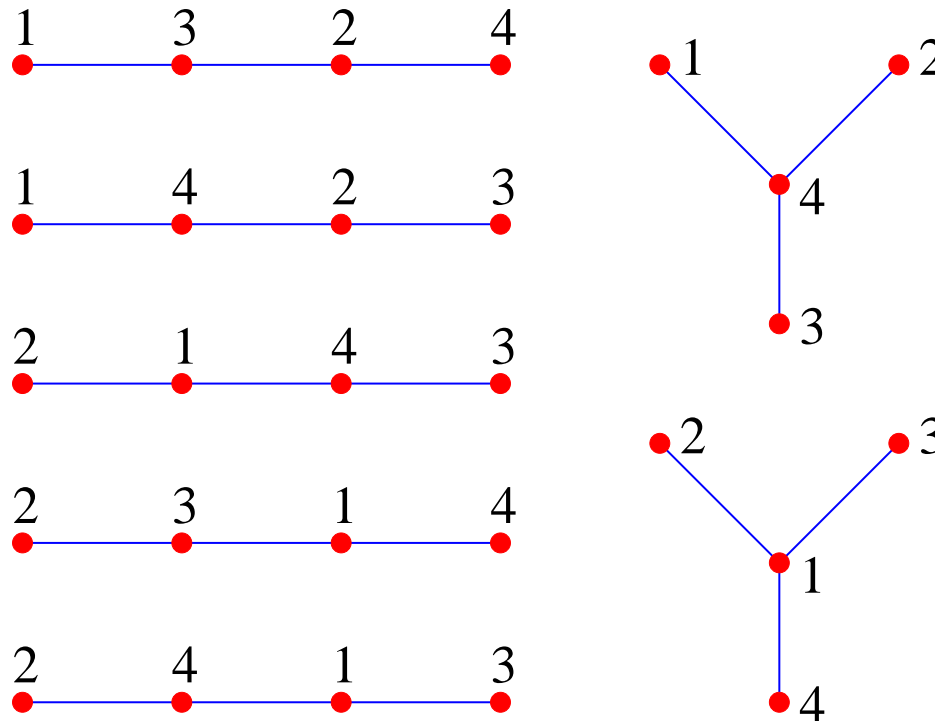
Athanasiadis: (difficult) proof using finite field method

Alternating trees

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An **alternating tree** on $[n]$ is a tree on the vertex set $[n]$ such that every vertex is either less than all its neighbors or greater than all its neighbors.



Alternating trees and \mathcal{L}_n

$f(n)$: number of alternating trees on $[n]$

Theorem (Kuznetsov, Pak, Postnikov, 1994).

$$f(n+1) = \frac{1}{2^n} \sum_{k=1}^n \binom{n}{k} (k+1)^{n-1}$$

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No combinatorial proof known!

The threshold arrangement

$$\mathcal{T}_n : x_i + x_j = 0, \quad 1 \leq i < j \leq n$$

$$\dim \mathcal{T}_n = n, \quad \text{rk } \mathcal{T}_n = n, \quad \#\mathcal{T}_n = \binom{n}{2}$$

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threshold graph:

- \emptyset is a threshold graph
- G threshold $\Rightarrow G \cup \{\text{vertex}\}$ threshold
- G threshold $\Rightarrow \text{join}(G, v)$ threshold

Char. poly. of threshold arrangements

Theorem. $r(\mathcal{T}_n) = \#$ threshold graphs on $[n]$.
Hence (by a known result on threshold graphs)

$$\sum_{n \geq 0} r(\mathcal{T}_n) \frac{x^n}{n!} = \frac{e^x(1-x)}{2-e^x}.$$

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Theorem.
$$\sum_{n \geq 0} \chi_{\mathcal{T}_n}(t) \frac{x^n}{n!} = (1+x)(2e^x - 1)^{(t-1)/2}$$

Small values of $\chi_{\mathcal{T}_n}(t)$

$$\chi_{\mathcal{T}_3}(t) = t^3 - 3t^2 + 3t - 1$$

$$\chi_{\mathcal{T}_4}(t) = t^4 - 6t^3 + 15t^2 - 17t + 7$$

$$\chi_{\mathcal{T}_5}(t) = t^5 - 10t^4 + 45t^3 - 105t^2 + 120t - 51.$$

Coefficients of $\chi_{\mathcal{T}_n}(t)$

Let

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Open: interpret a_i as the number of threshold graphs on $[n]$ with some property.

Minkowski space $\mathbb{R}^{1,3}$

$\mathbb{R}^{1,3}$: Minkowski spacetime with one time and three space dimensions

$$\boldsymbol{p} = (t, \boldsymbol{x}) \in \mathbb{R}^{1,3}, \quad \boldsymbol{x} = (x, y, z) \in \mathbb{R}^3$$

$$|\boldsymbol{p}|^2 = t^2 - |\boldsymbol{x}|^2 = t^2 - (x^2 + y^2 + z^2)$$

Ordering events in $\mathbb{R}^{1,3}$

Let $p_1, \dots, p_k \in \mathbb{R}^{1,3}$. In different reference frames (at constant velocities with respect to each other) these events can occur in different orders (but never violating causality).

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Let $p_1, \dots, p_k \in \mathbb{R}^{1,3}$. In different reference frames (at constant velocities with respect to each other) these events can occur in different orders (but never violating causality).

Main question: what is the maximum number of different orders in which these events can occur?

The hyperplane of simultaneity

Let $\mathbf{p}_1 = (t_1, \mathbf{x}_1)$, $\mathbf{p}_2 = (t_2, \mathbf{x}_2) \in \mathbb{R}^{1,3}$.

For a reference frame at velocity \mathbf{v} , the Lorentz transformation $\Rightarrow \mathbf{p}_1, \mathbf{p}_2$ occur at the same time if and only if

$$t_1 - t_2 = (\mathbf{x}_1 - \mathbf{x}_2) \cdot \mathbf{v}.$$

The set of all such $\mathbf{v} \in \mathbb{R}^3$ forms a hyperplane.

The Einstein arrangement

Thus the number of different orders in which the events can occur is the number of regions R of the **Einstein arrangement**

$$\mathcal{E} = \mathcal{E}(p_1, \dots, p_k)$$

defined by

$$t_i - t_j = (\mathbf{x}_1 - \mathbf{x}_2) \cdot \mathbf{v}, \quad 1 \leq i < j \leq k,$$

such that $|\mathbf{v}| < 1$ (the speed of light) for some $\mathbf{v} \in R$.

Intersection poset of \mathcal{E}

Can insure that $v \in R$ for all R by taking p_1, \dots, p_k sufficiently “far apart”.

Can maximize $r(\mathcal{E})$ for fixed k by choosing p_1, \dots, p_k generic.

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In this case, $L(\mathcal{E})$ is isomorphic to the **rank 3 truncation** of $L(\mathcal{B}_k) \cong \Pi_k$.

Coefficients of $\chi_{\mathcal{B}_k}(t)$

Recall

$$\begin{aligned}\chi_{\mathcal{B}_k}(t) &= t(t-1) \cdots (t-k+1) \\ &= c(k, k)t^k - c(k, k-1)t^{k-1} + \cdots ,\end{aligned}$$

where $c(k, i)$ is the number of permutations of $1, 2, \dots, k$ with i cycles (signless Stirling number of the first kind).

Computation of $r(\mathcal{E})$

Corollary.

$$\chi_{\mathcal{E}}(t) = c(k, k)t^3 - c(k, k-1)t^2 + c(k, k-2)t - c(k, k-3)$$

$$\Rightarrow r(\mathcal{E}) = c(k, k) + c(k, k-1) + c(k, k-2) + c(k, k-3)$$

$$= \frac{1}{48} (k^6 - 7k^5 + 23k^4 - 37k^3 + 48k^2 - 28k + 48)$$

